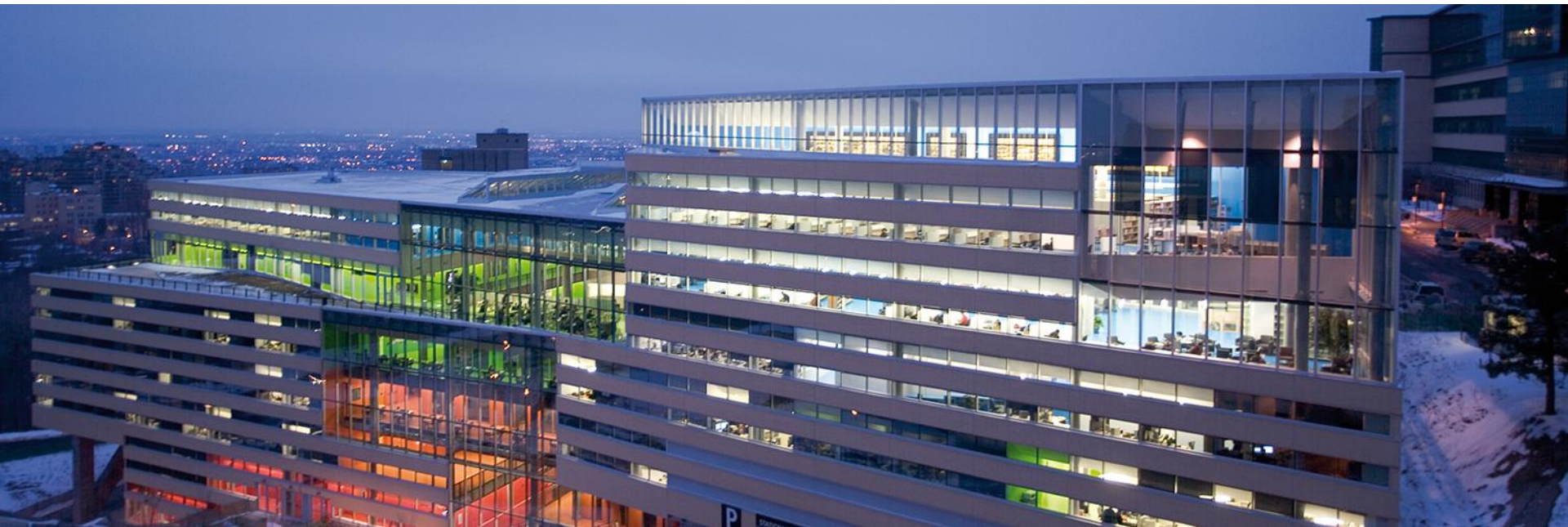# Modern Techniques For Query Evaluation on Highly Connected Datasets

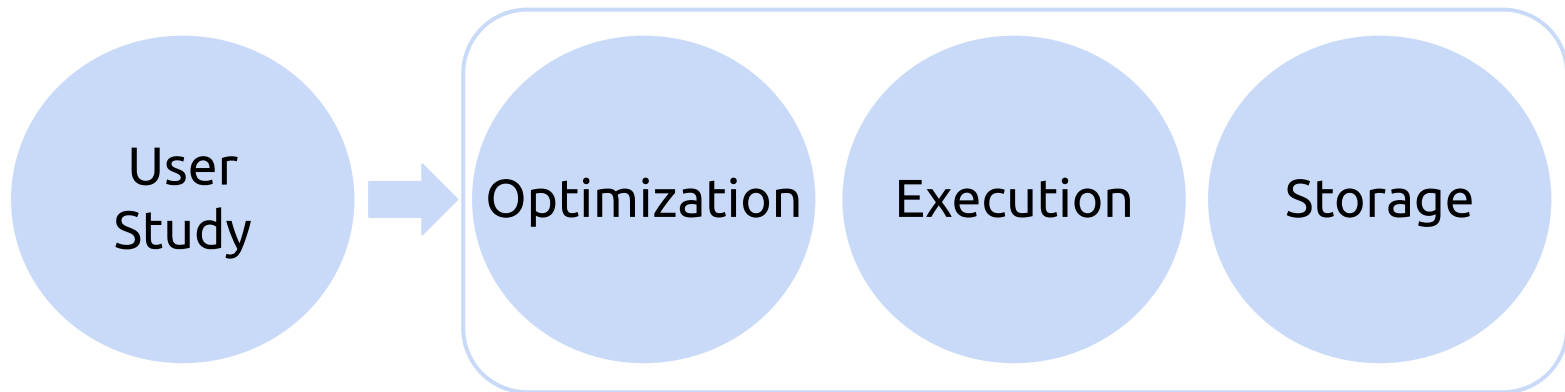## Amine Mhedhbi

Feb. 29th, 2024
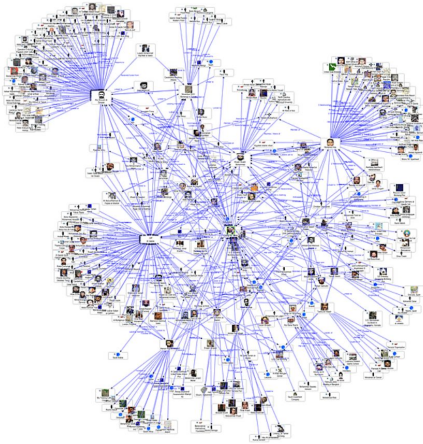
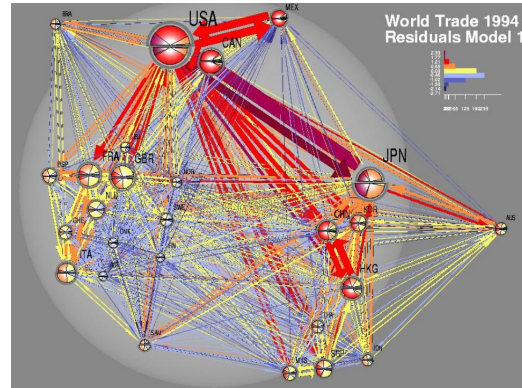I aim to design and implement data systems capable of *efficient graph data management.*
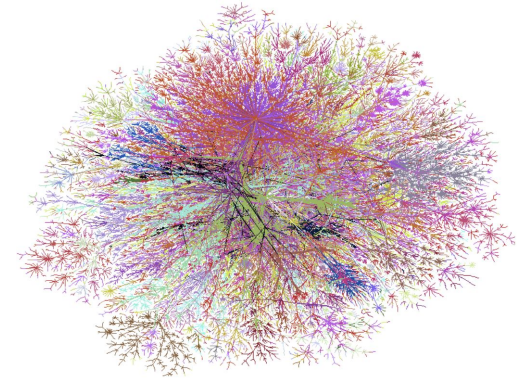
Modeling Application Data As Graph is Ubiquitous

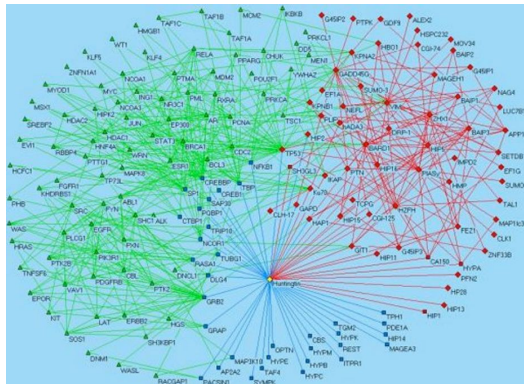# Modeling Application Data As Graph is Ubiquitous

Social Networks

Trade & Financial Connections

Internet / Web

Biological Networks

Linked Data *e.g.* Wiki

Road Networks

Relationships between entities are central to data analysis

&

Structure of the relationships provides insight

Biological Networks   Linked Data *e.g.* Wiki   Road Networks

Modeling Application Data As Graph is Ubiquitous

Social Networks
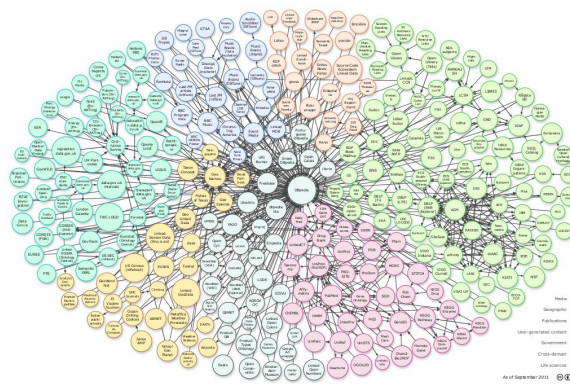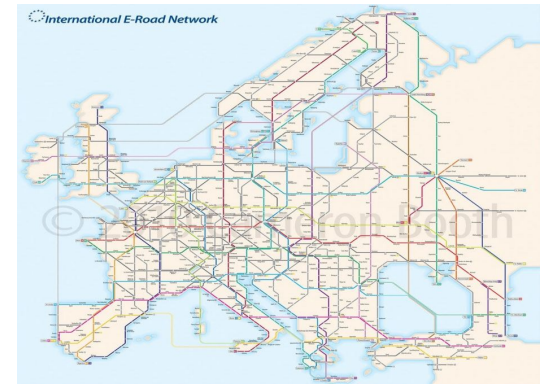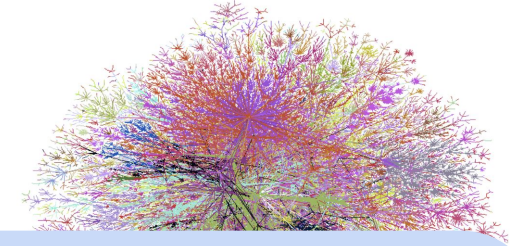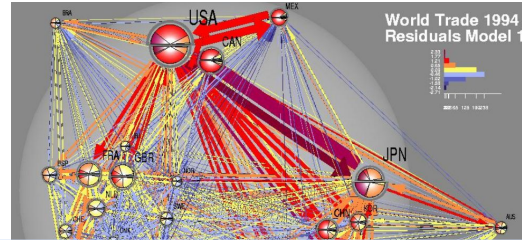

World Trade 1994
Residuals Model 1

Trade & Financial
Connections

Internet / Web

Biological Networks

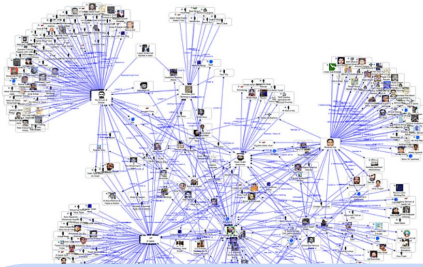Linked Data *e.g.* Wiki

International E-Road Network

Road Networks

Examples of Applications

Social Networks

Payment Services

Twitter Recommendation Example

Social Networks

Payment Services

WTF: The Who to Follow Service at Twitter. Gupta et al. WWW 2013.

Twitter Recommendation Example

Social Networks

Payment Services

WTF: The Who to Follow Service at Twitter. Gupta et al. WWW 2013.

Alipay Fraud Detection Example

Social Networks

Payment Services

Alipay Fraud Detection Example

Social Networks

Payment Services

Criminal   Merchant   Bank   loan

Real-time Constrained Cycle Detection in Large Dynamic Graphs. Qiu et al. VLDB 2018.

Alipay Fraud Detection Example

Social Networks

Payment Services

Real-time Constrained Cycle Detection in Large Dynamic Graphs. Qiu et al. VLDB 2018.
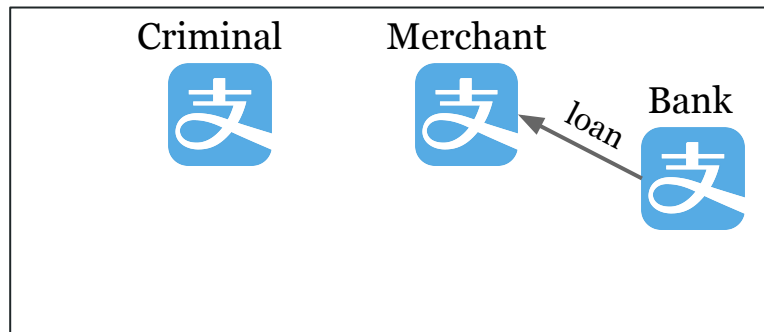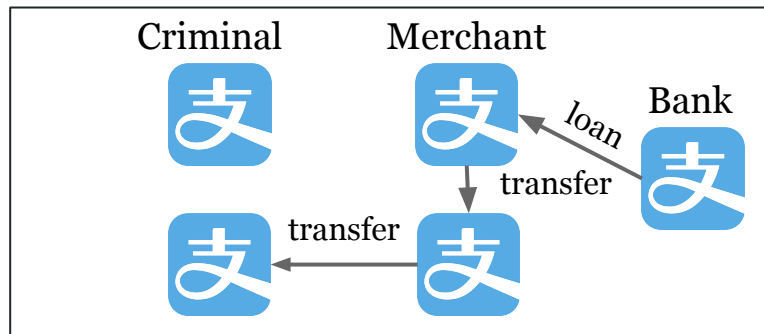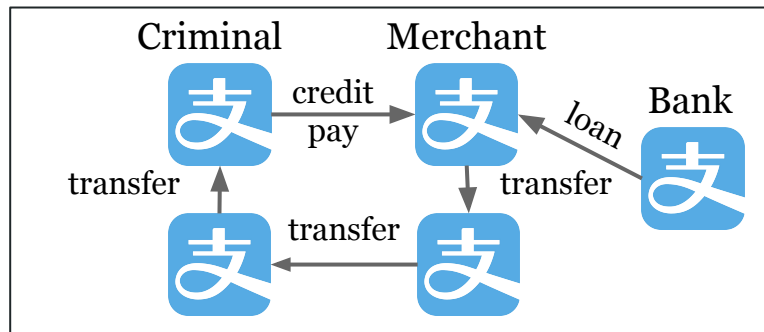
Alipay Fraud Detection Example

Social Networks

Payment Services

Real-time Constrained Cycle Detection in Large Dynamic Graphs. Qiu et al. VLDB 2018.

Examples of Applications

Social Networks

Payment Services

Graph Usage Study

The Ubiquity of Large Graphs and Surprising Challenges of Graph Processing

Siddhartha Sahu, Amine Mhedhbi, Semih Salihoglu, Jimmy Lin, M. Tamer Özsu
David R. Cheriton School of Computer Science
University of Waterloo
{s3sahu,amine.mhedhbi,semih.salihoglu,jimmylin,tamer.ozsu}@uwaterloo.ca

ABSTRACT

Graph processing is becoming increasingly prevalent across many application domains. In spite of this prevalence, there is little research about how graphs are actually used in practice. We conducted an online survey aimed at understanding: (i) the types of graphs users have; (ii) the graph computations users run; (iii) the types of graph software users use; and (iv) the major challenges users face when processing their graphs. We describe the participants' responses to our questions highlighting common patterns and challenges. We further reviewed user feedback in the mailing lists, bug reports, and feature requests in the source repositories of a large suite of software products for processing graphs. Through our review, we were able to answer some new questions that were raised by participants' responses and identify specific challenges that users face when using different classes of graph software. The participants' responses and data we obtained revealed surprising facts about graph processing in practice. In particular, real-world graphs represent a very diverse range of entities and are often very large, and scalability and visualization are undeniably the most pressing challenges faced by participants. We hope these findings can guide future research.

1. INTRODUCTION

Graph data representing connected entities and their relationships appear in many application domains, most naturally in social networks, the web, the semantic web, road maps, communication networks, biology, and finance, just to name a few examples. There has been a noticeable increase in the prevalence of work on graph processing both in research and in practice, evidenced by the surge in the number of different commercial and research software for managing and processing graphs. Examples include graph database systems [3,8,14,35,48,53], RDF engines [38,64,67], linear algebra software [6,46], visualization software [13,16], query languages [28,

52,55], and distributed graph processing systems [17,21,27]. In the academic literature, a large number of publications that study numerous topics related to graph processing regularly appear across a wide spectrum of research venues.

Despite their prevalence, there is little research on how graph data is actually used in practice and the major challenges facing users of graph data, both in industry and research. In April 2017, we conducted an online survey across 89 users of 22 different software products, with the goal of answering 4 high-level questions:

(i) What types of graph data do users have?
(ii) What computations do users run on their graphs?
(iii) Which software do users use to perform their computations?
(iv) What are the major challenges users face when processing their graph data?

Our major findings are as follows:

- *Variety:* Graphs in practice represent a very wide variety of entities, many of which are not naturally thought of as vertices and edges. Most surprisingly, traditional enterprise data comprised of products, orders, and transactions, which are typically seen as the perfect fit for relational systems, appear to be a very common form of data represented in participants' graphs.
- *Ubiquity of Very Large Graphs:* Many graphs in practice are very large, often containing over a billion edges. These large graphs represent a very wide range of entities and belong to organizations at all scales from very small enterprises to very large ones. This refutes the sometimes heard assumption that large graphs are a problem for only a few large organizations such as Google, Facebook, and Twitter.
- *Challenge of Scalability:* Scalability is unequivocally the most pressing challenge faced by participants. The ability to process very large graphs efficiently seems to be the biggest limitation of existing software.
- *Visualization:* Visualization is a very popular and central task in participants' graph processing pipelines. After scalability, participants indicated visualization as their second most pressing challenge, tied with challenges in graph query languages.
- *Prevalence of RDBMSes:* Relational databases still play an important role in managing and processing graphs.

The ubiquity of large graphs and surprising challenges of graph processing: extended survey

Siddhartha Sahu[1] · Amine Mhedhbi[1] · Semih Salihoglu[1] · Jimmy Lin[1] · M. Tamer Özsu[1]

Abstract

Graph processing is becoming increasingly prevalent across many application domains. In spite of this prevalence, there is little research about how graphs are actually used in practice. We performed an extensive study that consisted of an online survey of 89 users, a review of the mailing lists, source repositories, and white papers of a large suite of graph software products, and in-person interviews with 6 users and 2 developers of these products. Our online survey aimed at understanding: (i) the types of graphs users have; (ii) the graph computations users run; (iii) the types of graph software users use; and (iv) the major challenges users face when processing their graphs. We describe the participants' responses to our questions highlighting common patterns and challenges. Based on our interviews and survey of the rest of our sources, we were able to answer some new questions that were raised by participants' responses to our online survey and understand the specific applications that use graph data and software. Our study revealed surprising facts about graph processing in practice. In particular, real-world graphs represent a very diverse range of entities and are often very large, scalability and visualization are undeniably the most pressing challenges faced by participants, and data integration, recommendations, and fraud detection are very popular applications supported by existing graph software. We hope these findings can guide future research.

1 Introduction

Graph data representing connected entities and their relationships appear in many application domains, most naturally in social networks, the Web, the Semantic Web, road maps, communication networks, biology, and finance, just to name a few examples. There has been a noticeable increase in the

prevalence of work on graph processing both in research and in practice, evidenced by the surge in the number of different commercial and research software for managing and processing graphs. Examples include graph database systems [13,20,26,49,65,73,90], RDF engines [52,96], linear algebra software [17,63], visualization software [25,29], query languages [41,72,78], and distributed graph processing systems [30,34,40]. In the academic literature, a large number of publications that study numerous topics related to graph processing regularly appear across a wide spectrum of research venues.

Despite their prevalence, there is little research on how graph data are actually used in practice and the major challenges facing users of graph data, both in industry and in research. In April 2017, we conducted an online survey across 89 users of 22 different software products, with the goal of answering 4 high-level questions:

Objectives

What kind of graph data, computations, software, and major challenges **industry users** have?

Objectives

What kind of graph data, computations, software, and major challenges **industry users** have?

Some Major Findings

1. Graphs are very large!
2. Scalability is the most pressing challenge!
3. ML on graphs is very popular (> 85% of respondents have ML workloads)!

Apps Store Facts or Events

Property Graph Data Model

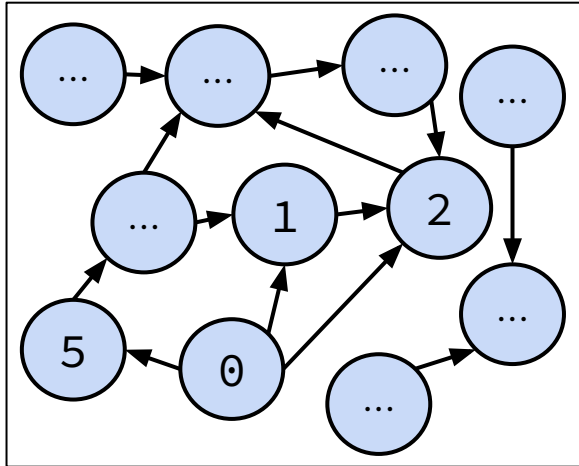| (FROM, TO) |
|:---:|
| (0, 1) |
| (0, 2) |
| (0, 5) |
| ... |

Relational Model

Property Graph Data Model
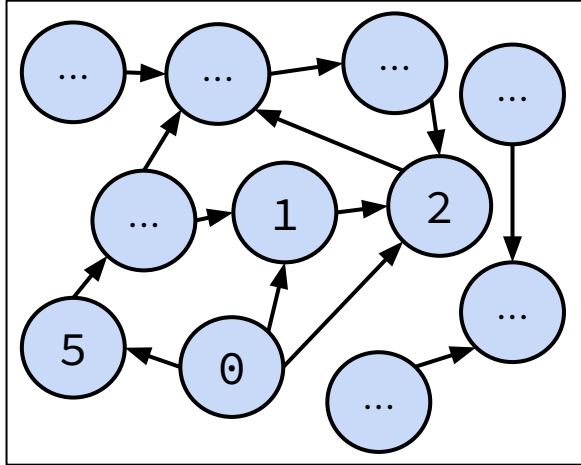
(FROM, TO)
(0, 1)
(0, 2)
(0, 5)
...

Relational Model

Relational algebra primitives!

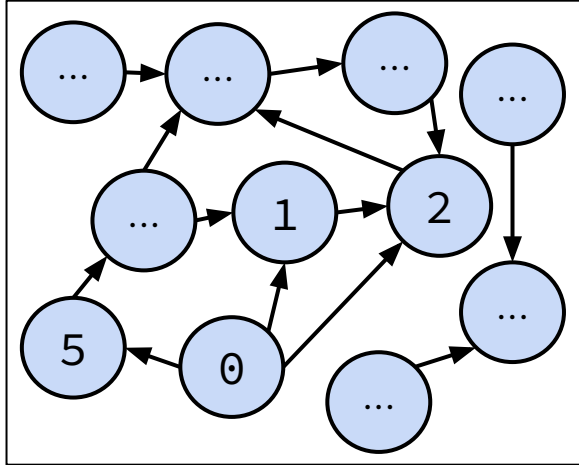→ Emphasis on workload i.e., queries and dataset characteristics.

Highly Connected Dataset

1. <u>Highly Connected data</u>:

   Lots of many-to-many relationships (N-to-M cardinality) !!

Highly Connected Dataset

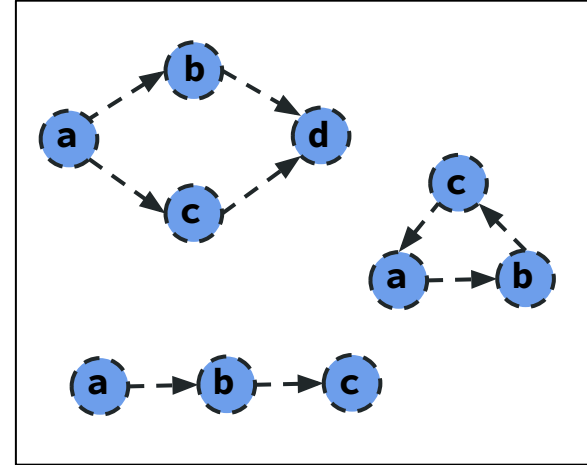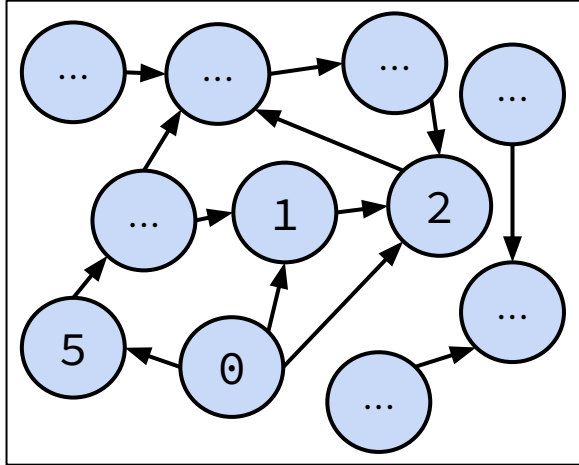**Structure-based Queries**

1. <u>Highly Connected data</u>:

   Lots of many-to-many relationships (N-to-M cardinality) !!

Workload Characteristics

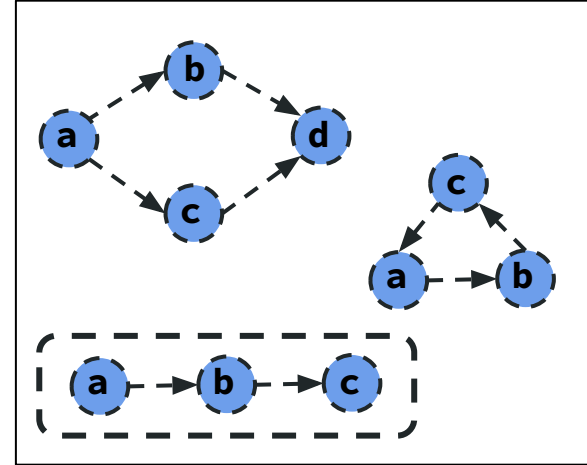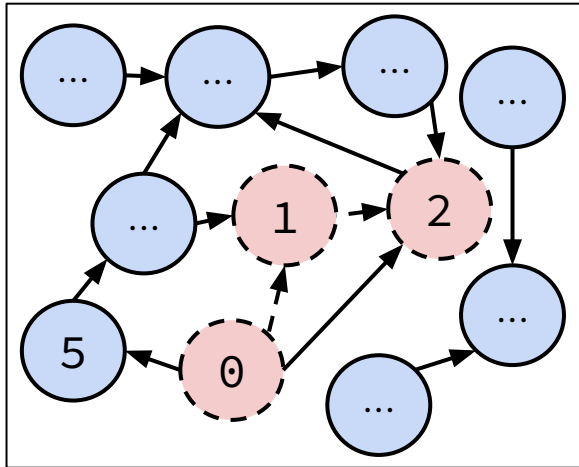**Highly Connected Dataset**

Structure-based Queries
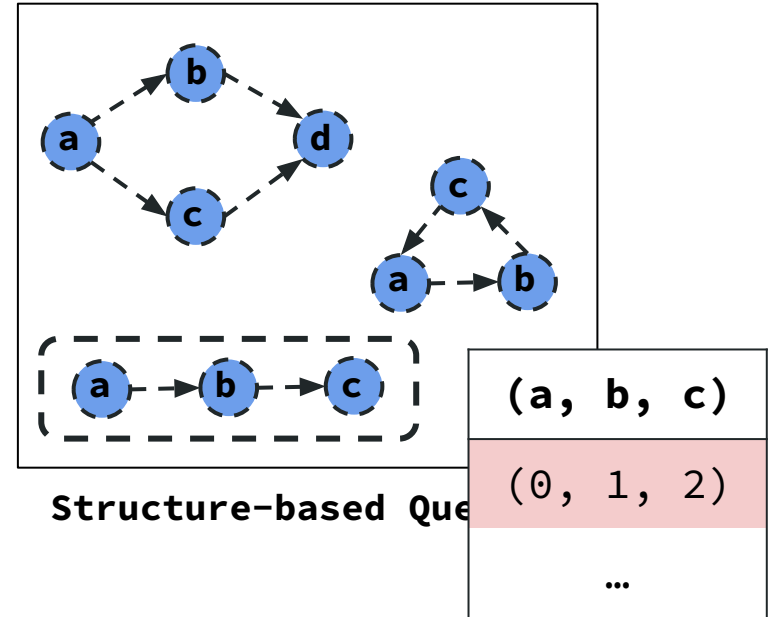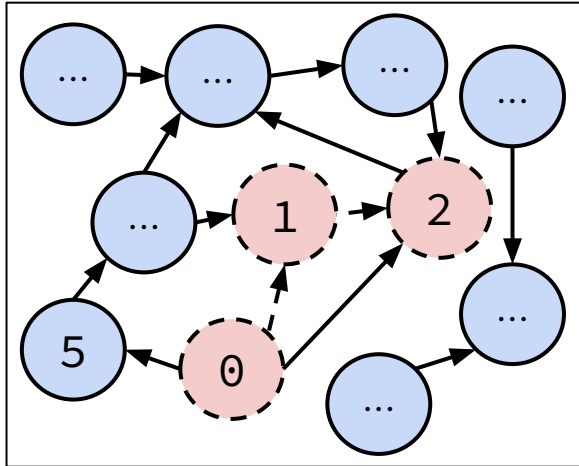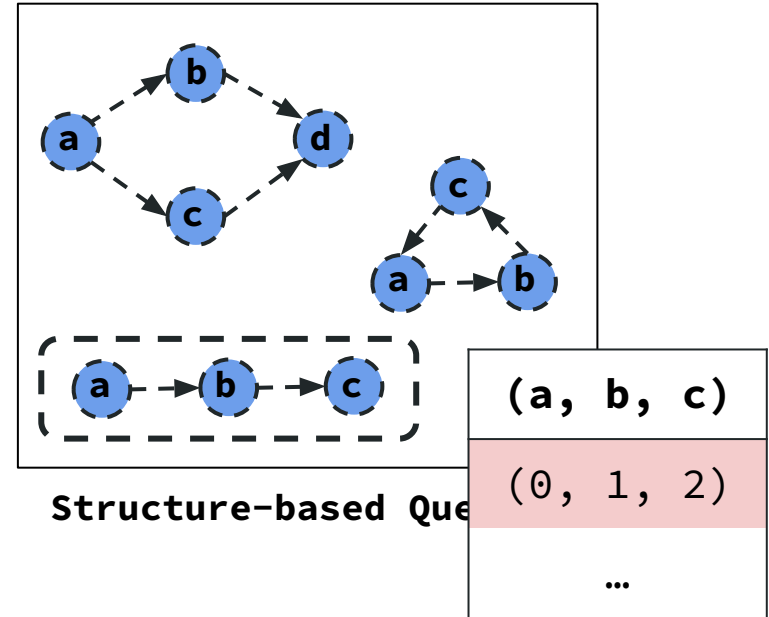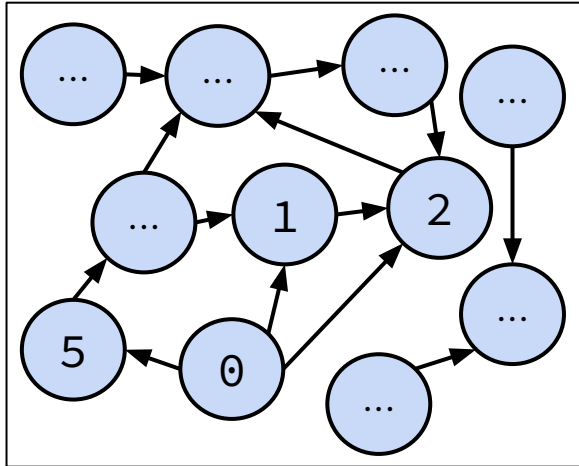
1. Highly Connected data:

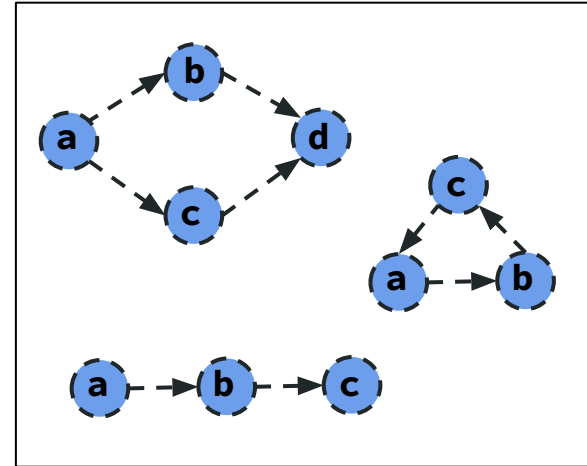   Lots of many-to-many relationships (N-to-M cardinality) !!

Highly Connected Dataset

Structure-based Que...

(a, b, c)
(0, 1, 2)
…

1. <u>Highly Connected data</u>:

    Lots of many-to-many relationships (N-to-M cardinality) !!

Highly Connected Dataset

Structure-based Queries

(a, b, c)
(0, 1, 2)
…

1. <u>Highly Connected data</u>:

    Lots of many-to-many relationships (N-to-M cardinality) !!

2. <u>Structure-based Queries</u>:

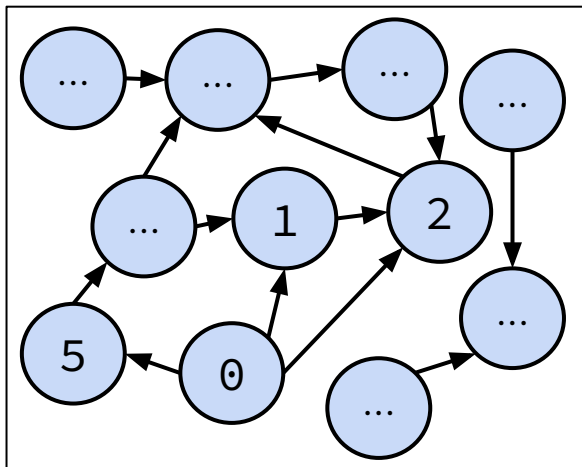   Queries enumerate graph patterns (*complex many-to-many joins*)
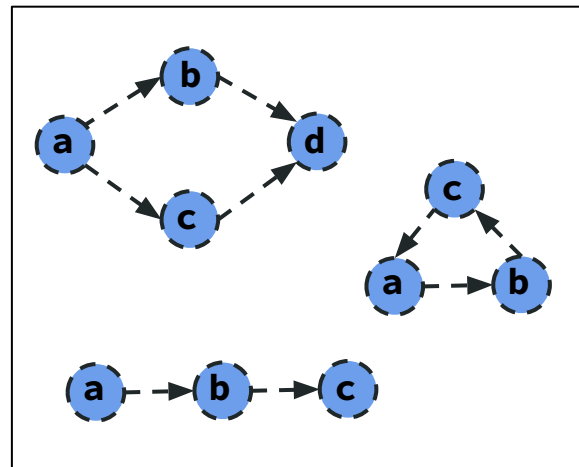
Highly Connected Dataset

Structure-based Queries

Colloquially called "graph workloads" aka Querying

Workload Characteristics

Highly Connected Dataset

Structure-based Queries

Colloquially called "graph workloads" aka Querying

Interest in **performance issues.**

Challenging workload due to complex many-to-many joins
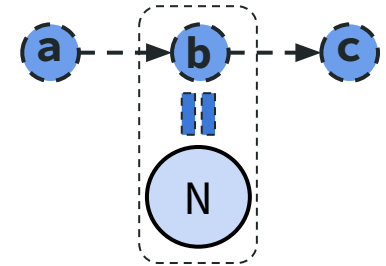
→ Existing data systems come short.

Workload Challenges - Many-to-many Joins

On a financial network, find the accounts whose transactions
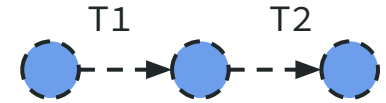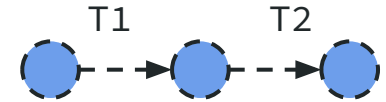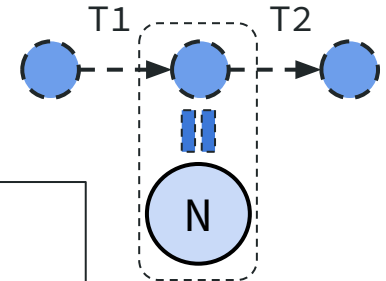are facilitated by user with account ID N?

Transactions

src	dst	amount	date

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

On a financial network, find the accounts whose transactions
are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
```

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
  FROM Transactions T1, Transactions T2
WHERE T1.dst = T2.src
```

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```
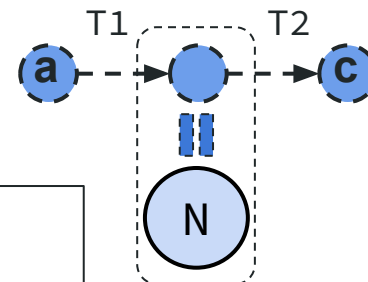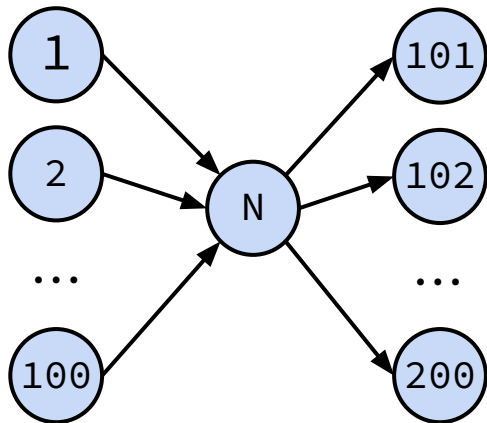
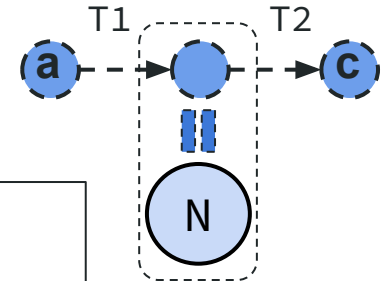On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```
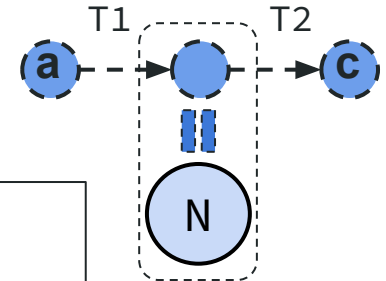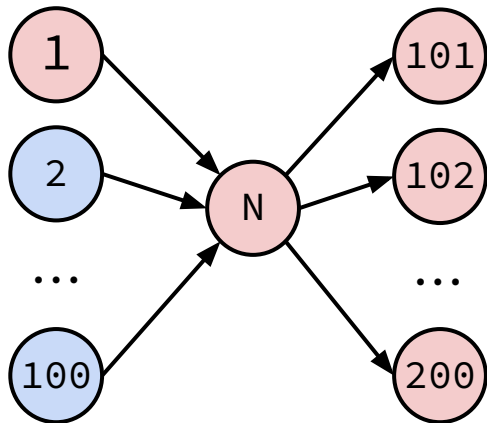
T1   T2

N

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```
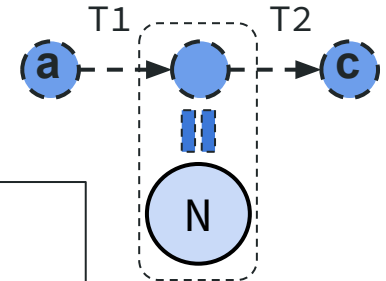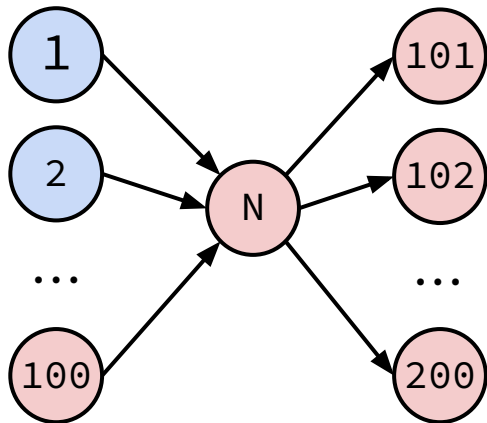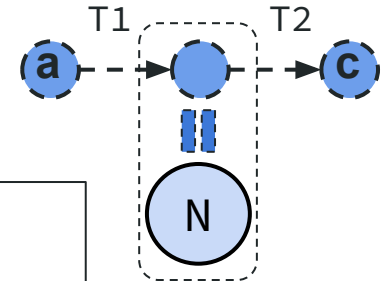
T1 T2

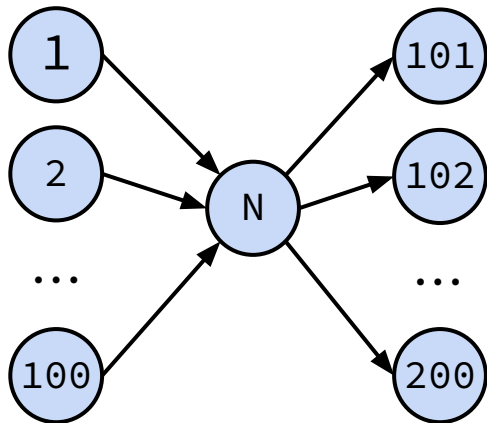a → N → c

Workload Challenges - Many-to-many Joins

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
   FROM Transactions T1, Transactions T2
  WHERE T1.dst = T2.src
    AND T1.dst = N
```

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```
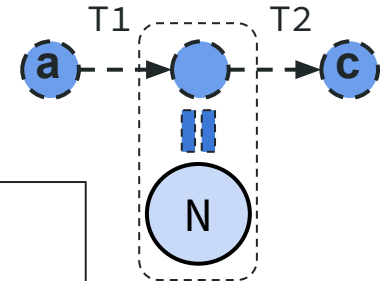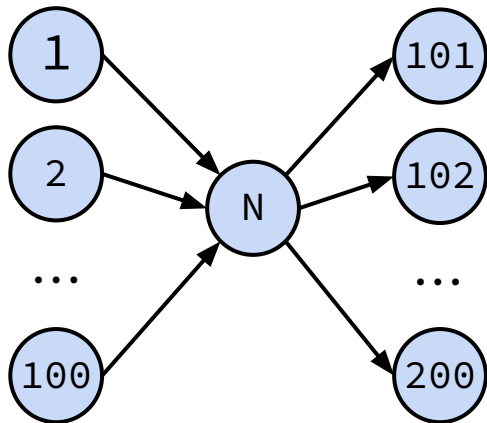
(a, c)
(1, 101)
(1, ...)
(1, 200)

100 tuples

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```

(a, c)
(1, 101)
(1, ...)
(1, 200)
...
(100, 101)
(100, ...)
(100, 200)

100 tuples

39

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```

(a, c)
(1, 101)
(1, ...)
(1, 200)
...
(100, 101)
(100, ...)
(100, 200)

200 edges → 10,000 !!!

On a financial network, find the accounts whose transactions are facilitated by user with account ID N?

Transactions

src	dst	amount	date

```
SELECT T1.src as a, T2.dst as c
  FROM Transactions T1, Transactions T2
 WHERE T1.dst = T2.src
   AND T1.dst = N
```

(a, c)
(1, 101)
(1, ...)
(1, 200)
...
(100, 101)
(100, ...)
(100, 200)

200 edges → 10,000 !!!

Explosion in intermediate results size!!!

DBMSs Evaluating Graph Workloads

Analytical RDBMSs do not
optimize for graph workloads

DBMSs Evaluating Graph Workloads

Analytical RDBMSs do not optimize for graph workloads

New graph DBMSs (GDBMSs) optimize for graph workloads

DBMSs Evaluating Graph Workloads

Analytical RDBMSs do not optimize for graph workloads

New graph DBMSs (GDBMSs) optimize for graph workloads

High-level Research Question – *How should Database Management Systems (DBMSs) be architected to optimize for analytical graph workloads?*

Revisiting DBMS Components

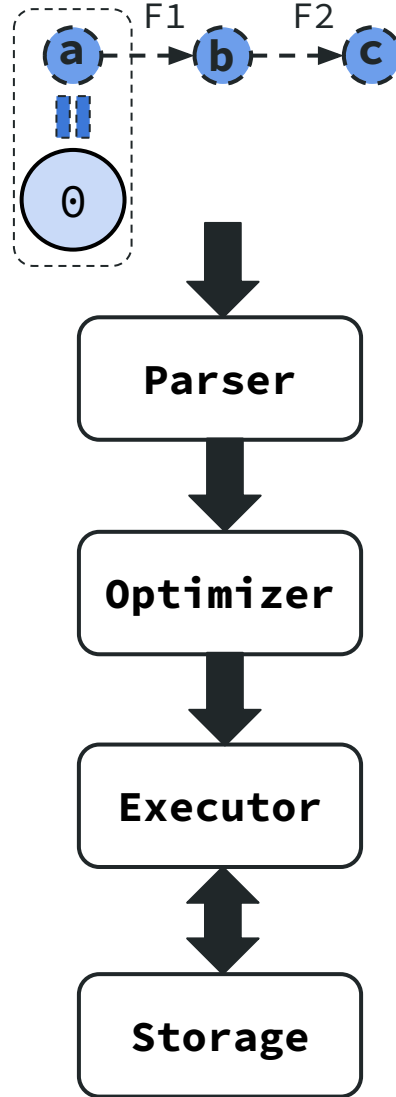*Query in Declarative Language e.g., SQL, Cypher, etc.*

```
Parser
  ↓
Optimizer
  ↓
Executor
  ↕
Storage
```
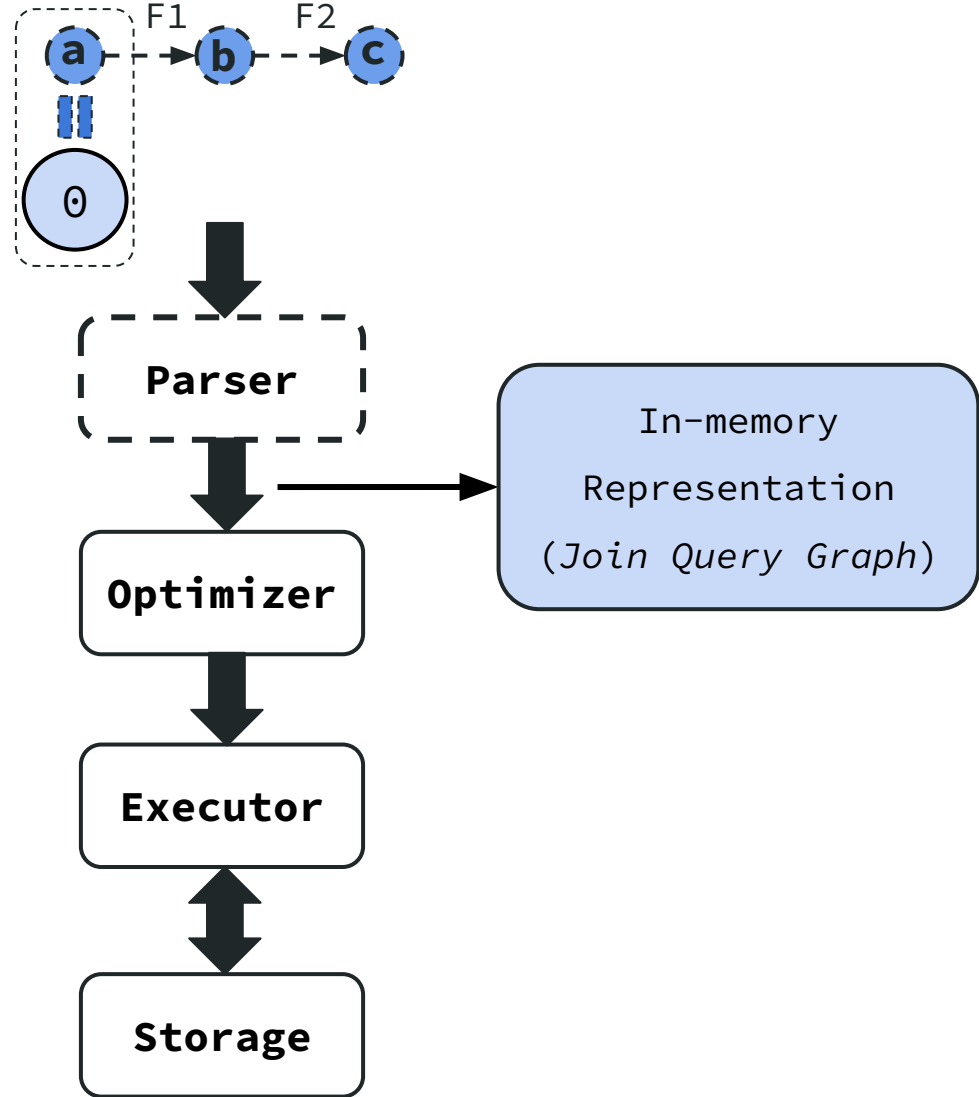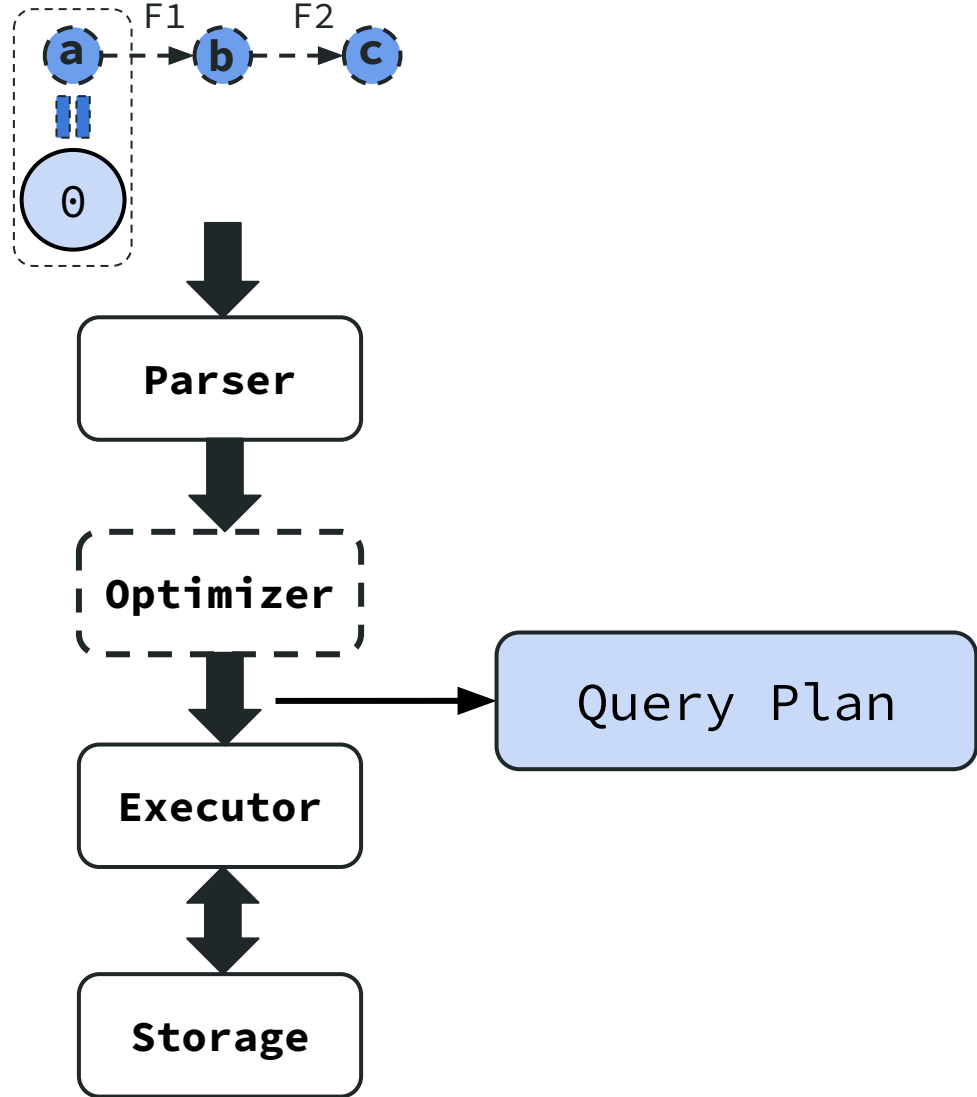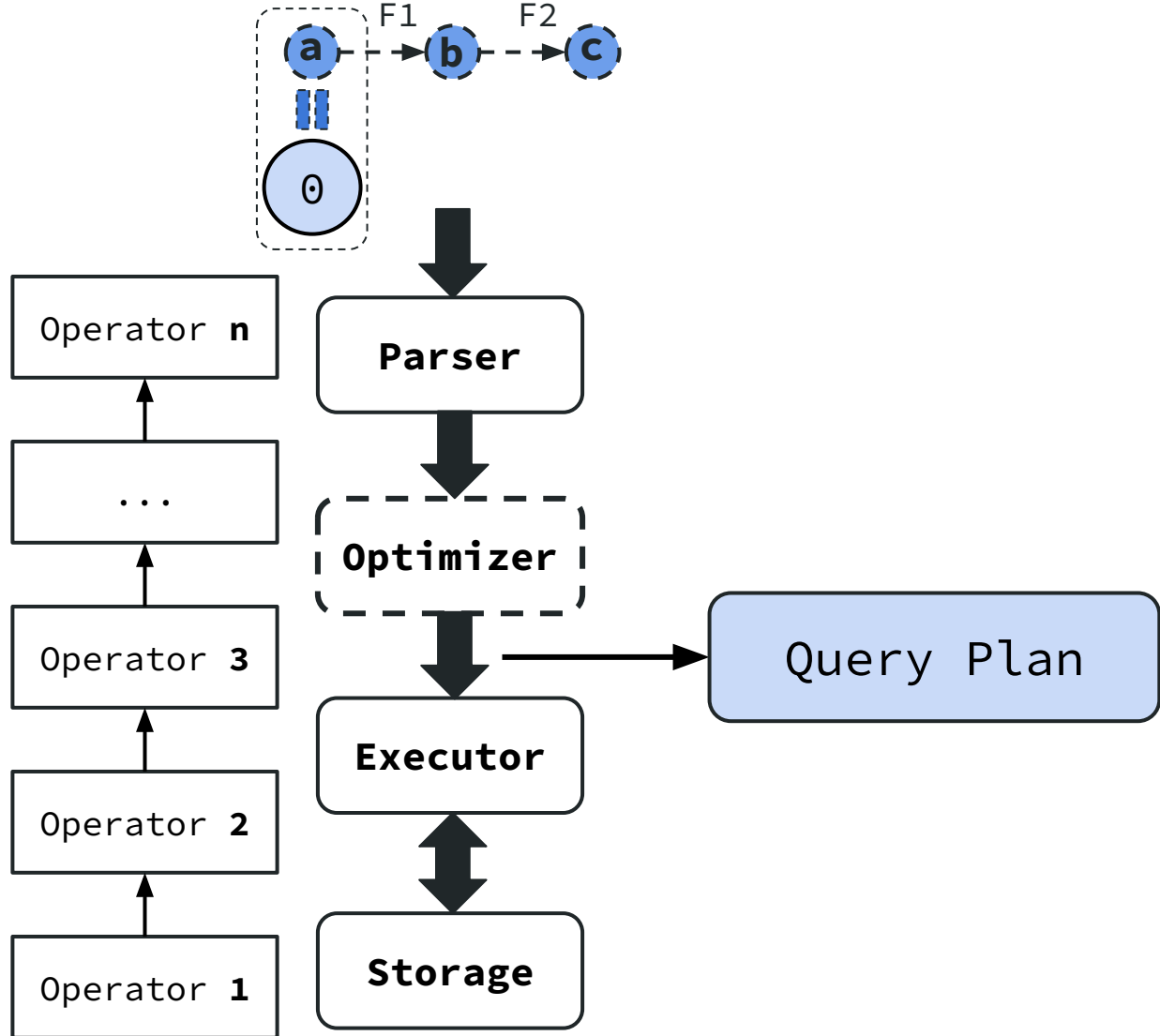
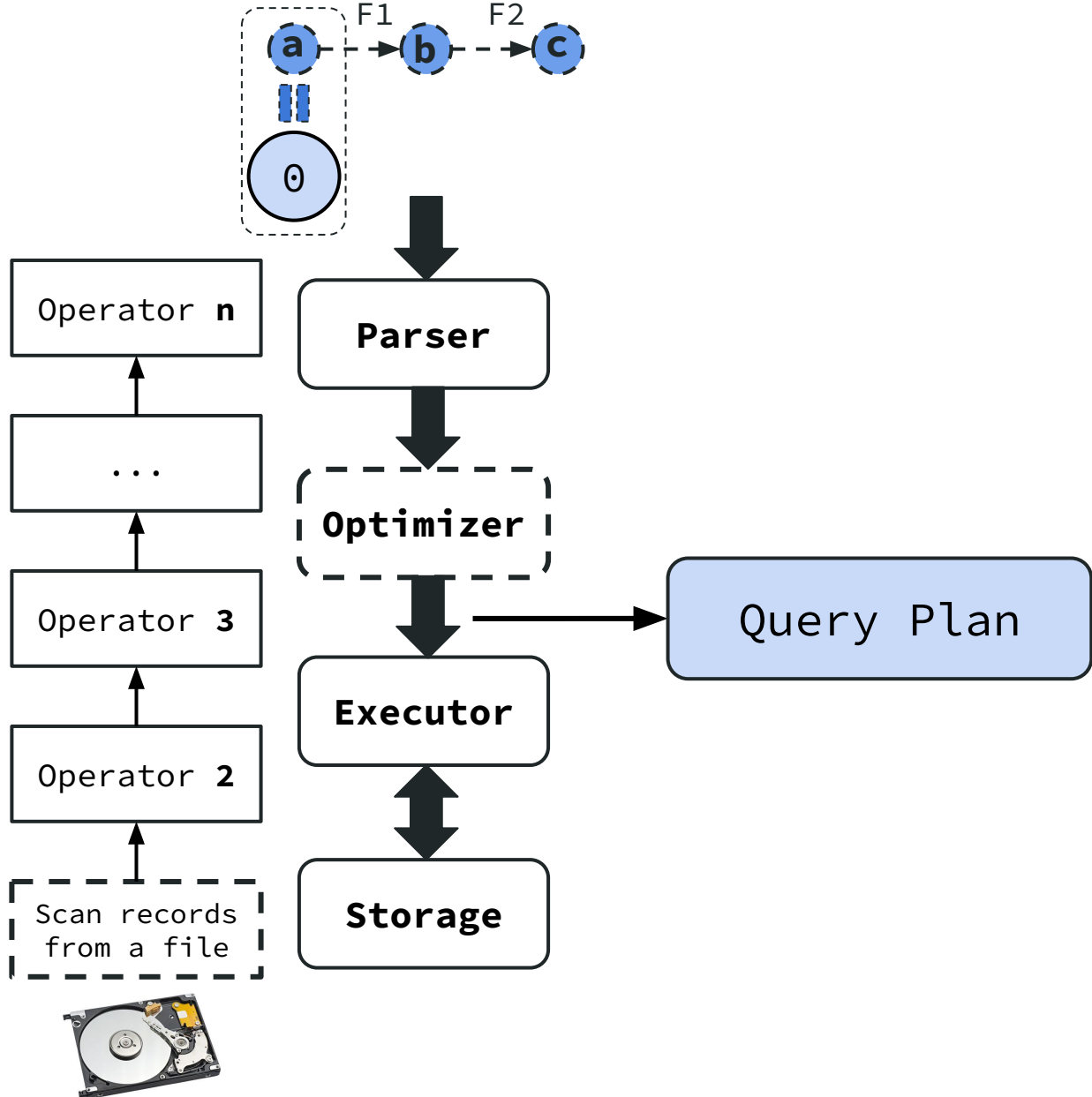# Revisiting DBMS Components (Example)

Revisiting DBMS Components (Example)

Cypher
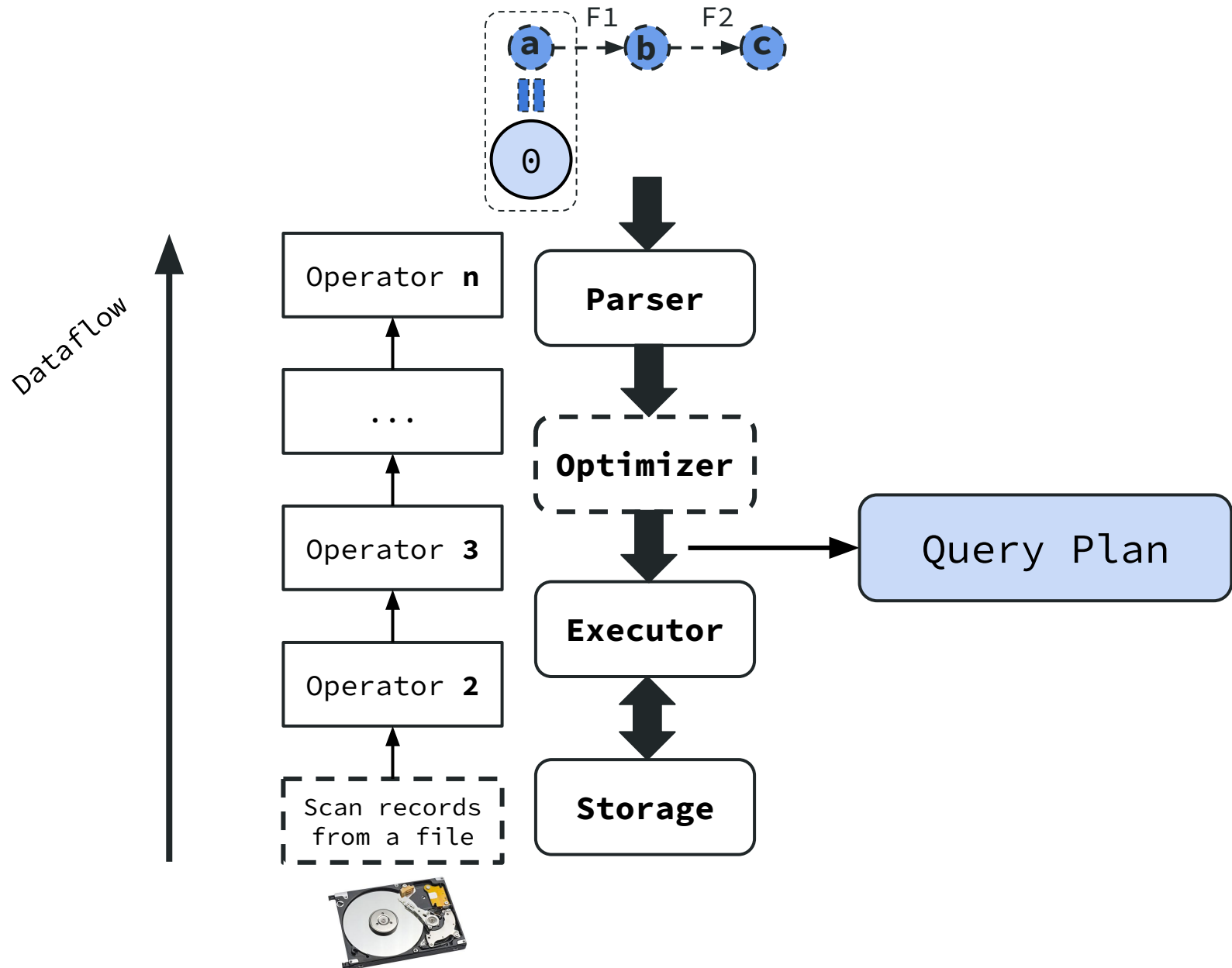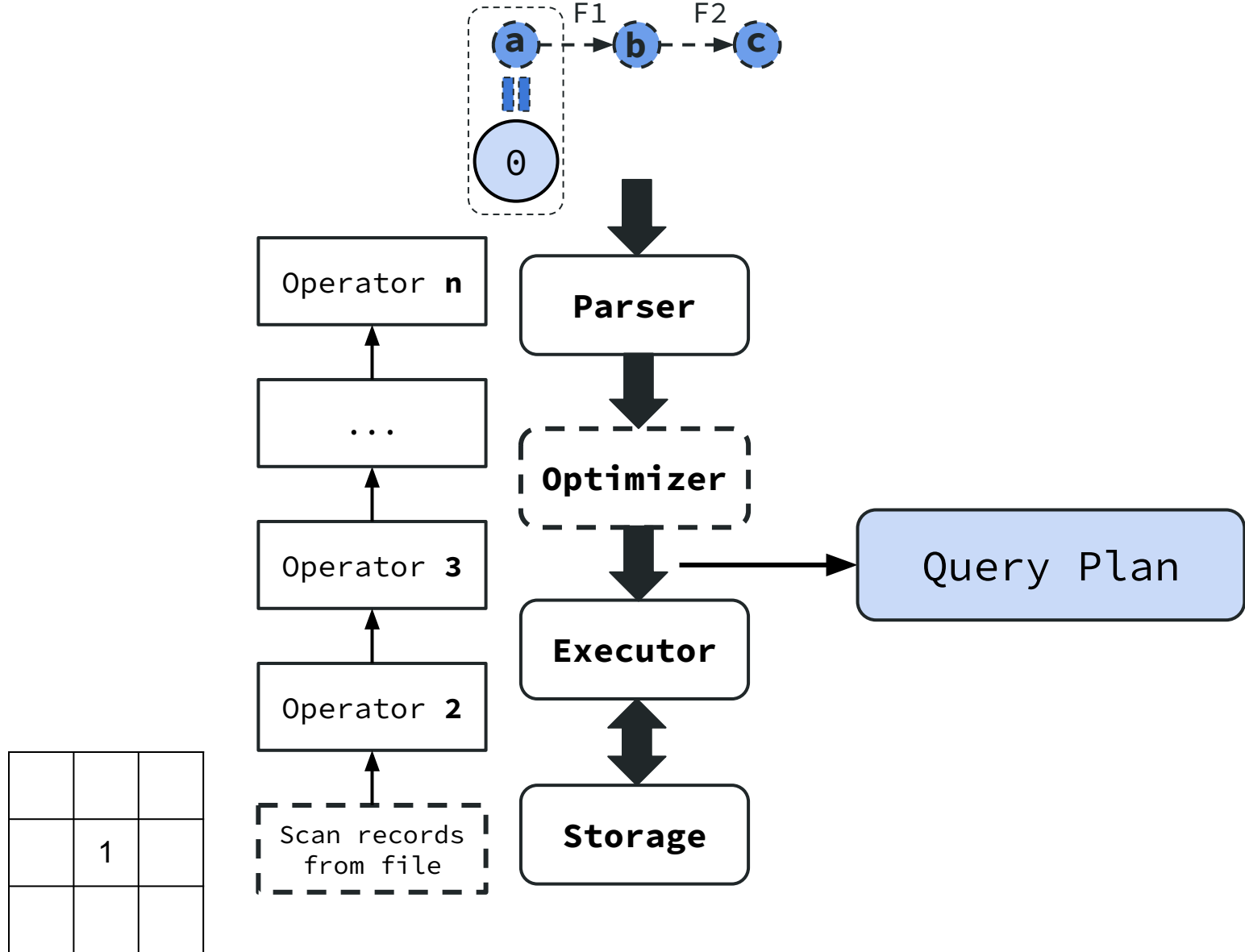

F1   F2

a - - → b - - → c

Parser

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Cypher

```
MATCH (a)-[Follows]->(b),
      (b)-[Follows]->(c)
```


Parser

Optimizer

Executor

Storage

Cypher

```
MATCH (a)-[Follows]->(b),
      (b)-[Follows]->(c)
WHERE  a.ID = 0
```


Parser

Optimizer

Executor

Storage

Cypher

```
MATCH  (a)-[Follows]->(b),
       (b)-[Follows]->(c)
WHERE  a.ID = 0
RETURN b.ID, c.ID
```

F1   F2

a -- -> b -- -> c

0

Parser

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Cypher

```
MATCH (a)-[Follows]->(b),
      (b)-[Follows]->(c)
WHERE  a.ID = 0
RETURN b.ID, c.ID
```

SQL

Parser

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Cypher

```
MATCH  (a)-[Follows]->(b),
       (b)-[Follows]->(c)
WHERE  a.ID = 0
RETURN b.ID, c.ID
```

SQL

```
FROM   Follows F1,
       Follows F2
```

Parser

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Cypher

```
MATCH  (a)-[Follows]->(b),
       (b)-[Follows]->(c)
WHERE   a.ID = 0
RETURN b.ID, c.ID
```

SQL

```
FROM    Follows F1,
        Follows F2
WHERE   F1.FROM = 0
  AND   F1.TO = F2.FROM
```


Parser

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Cypher

```
MATCH (a)-[Follows]->(b),
      (b)-[Follows]->(c)
WHERE  a.ID = 0
RETURN b.ID, c.ID
```

SQL

```
SELECT F2.FROM as b,
       F2.TO   as c
FROM   Follows F1,
       Follows F2
WHERE  F1.FROM = 0
  AND  F1.TO = F2.FROM
```

Parser

Optimizer

Executor

Storage

F1   F2

a --→ b --→ c

0

Parser

→ In-memory Representation (*Join Query Graph*)

Optimizer

Executor

Storage

Revisiting DBMS Components (Example)

Operator **n**

...

Operator **3**

Operator **2**

Scan records
from a file

Parser

Optimizer

Executor

Storage

Query Plan

F1   F2

a → b → c

0

DataFlow

Operator **n**

...

Operator **3**

Operator **2**

Scan records from a file

Parser

Optimizer

Executor

Storage

Query Plan

F1   F2

a → b → c

0

Operator **n**

...

Operator **3**

Operator **2**

Scan records from file

Parser

Optimizer

Executor

Storage

Query Plan

1

F1   F2

a → b → c

0

| Operator **n** |
| ... |
| Operator **3** |
| Operator **2** |
| Scan records from file |

	2	

Parser

Optimizer

Executor

Storage

Query Plan

F1   F2

a → b → c

0

Operator **n**

...

Operator **3**

Operator **2**

Scan records from file

	3	

Parser

Optimizer

Executor

Storage

Query Plan

User Study → Optimization Execution Storage

Novel Join Algorithms
VLDB 2019, TODS 2021
Compressed Representations
VLDB 2021

Query Processing Techniques Overview

Insights

1. Use novel join algorithms to remove unnecessary intermediate results!

2. Use compression to reduce the size of necessary intermediate results!

Query Processing Techniques Overview

Insights

1. Use novel join algorithms to remove unnecessary intermediate results!

> *Traditional Joins are*
> *suboptimal!*
>
> *Novel Join Algorithms*
> *correct the suboptimality*

Cyclic Pattern/Join Query

2. Use compression to reduce the size of necessary intermediate results!

Query Processing Techniques Overview

Insights

1. Use novel join algorithms to remove unnecessary intermediate results!

Traditional Joins are suboptimal!

Novel Join Algorithms correct the suboptimality

Cyclic Pattern/Join Query

2. Use compression to reduce the size of necessary intermediate results!

Large results part of the final output

Contain a lot of redundancy → compression

Acyclic Pattern/Join Query

Query Processing Techniques Overview

Insights

1. Use novel join algorithms to remove unnecessary intermediate results!

2. Use compression to reduce the size of necessary intermediate results!

Query Processing Techniques Overview

> ### **<u>Insights</u>**
>
> 1. Use novel join algorithms to remove unnecessary intermediate results!
>
> 2. Use compression to reduce the size of necessary intermediate results!

<u>Our approach provides</u>:

1. Up to ~10-70x speedups over State-of-the-art!

2. Queries run to completion!

Query Processing Techniques Overview

> ## <u>Insights</u>
>
> 1. Use novel join algorithms to remove unnecessary intermediate results!
>
> 2. Use compression to reduce the size of necessary intermediate results!

<u>Our approach provides</u>:

1. **Up to ~10-70x speedups over State-of-the-art!**

2. **Queries run to completion!**

- **Novel Join Algorithms**

  → Worst-case Optimal Joins

- **Compressed Representations**

  → Factorized Representations

- **Novel Join Algorithms**
  → Worst-case Optimal Joins

- **Compressed Representations**
  → Factorized Representations

Traditional Binary Joins Example

Query

Dataset

Query

Dataset

(a, b)	⋈	(b, c)	⋈	(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Query

Dataset

(a, b)	⋈	(b, c)	⋈	(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

78

Query

Dataset

(a, b)		**(b, c)**		**(a, c)**
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Query

(a, b, c)

(0, n, n+1)

...

(n−1, n, 2n)

(a, b)	⋈	**(b, c)**	⋈	**(a, c)**
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Dataset

Traditional Binary Joins Example

Query

(a, b, c)

(a, b, c)
(0, n, n+1)
...
(n−1, n, 2n)

(a, b)		(b, c)		(a, c)
(0, n)		(0, n)		(0, n)
...	⋈	...	⋈	...
(n, 2n)		(n, 2n)		(n, 2n)

Dataset

N

Query

Dataset

N^2 tuples

$$(a, b, c)$$

$$(0, n, n+1)$$

$$...$$

$$(n-1, n, 2n)$$

$$(a, b) \quad \bowtie \quad (b, c) \quad \bowtie \quad (a, c)$$

$$(0, n) \qquad (0, n) \qquad (0, n)$$

$$... \qquad ... \qquad ...$$

$$(n, 2n) \qquad (n, 2n) \qquad (n, 2n)$$

Query

(a, b, c)

(0, n, n+1)

...

(n−1, n, 2n)

$\}$ N^2 tuples

Dataset

(a, b) ⋈ **(b, c)** ⋈ **(a, c)**

(0, n)

...

(n, 2n)

(0, n)

...

(n, 2n)

(0, n)

...

(n, 2n)

Query

Dataset

(a, b, c)

(0, n, n+1)

(a, b, c)

(0, n, n+1)

...

(n−1, n, 2n)

N^2

(a, b) ⋈ **(b, c)** ⋈ **(a, c)**

(a, b)	(b, c)	(a, c)
(0, n)	(0, n)	(0, n)
...	...	...
(n, 2n)	(n, 2n)	(n, 2n)

Problem With Binary Join Plans
Prohibitively large # intermediate results for some Q.

For triangles, **AGM bound** is $N^{3/2}$
[FOCS 2008, PODS 2012]

Dataset

(a, b, c)

(0, n, n+1)

(a, b, c)

(0, n, n+1)

...

(n-1, n, 2n)

N^2

⋈

(a, b)

(0, n)

...

(n, 2n)

⋈

(b, c)

(0, n)

...

(n, 2n)

(a, c)

(0, n)

...

(n, 2n)

Problem With Binary Join Plans
Prohibitively large # intermediate results for some Q.

For triangles, **AGM bound** is $N^{3/2}$
[FOCS 2008, PODS 2012]

→ *Worst-case optimal joins correct for the suboptimality (Generic Join)*
SIGMOD Record *[Ngo et al. 2013]*

Dataset

Worst-case Optimal Joins (WCOJs) Example

Query

Dataset

Query

Dataset

(a, b)	⋈	(b, c)	⋈	(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Join Attribute Ordering (JAO)

[a,b,c]

1. Find set of a's
2. Given an (a), find b's
3. Given an (a,b), find c's

Query

Dataset

(a, b)	⋈	(b, c)	⋈	(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

90

Worst-case Optimal Joins (WCOJs) Example

Join Attribute Ordering (JAO)

[a,b,c]
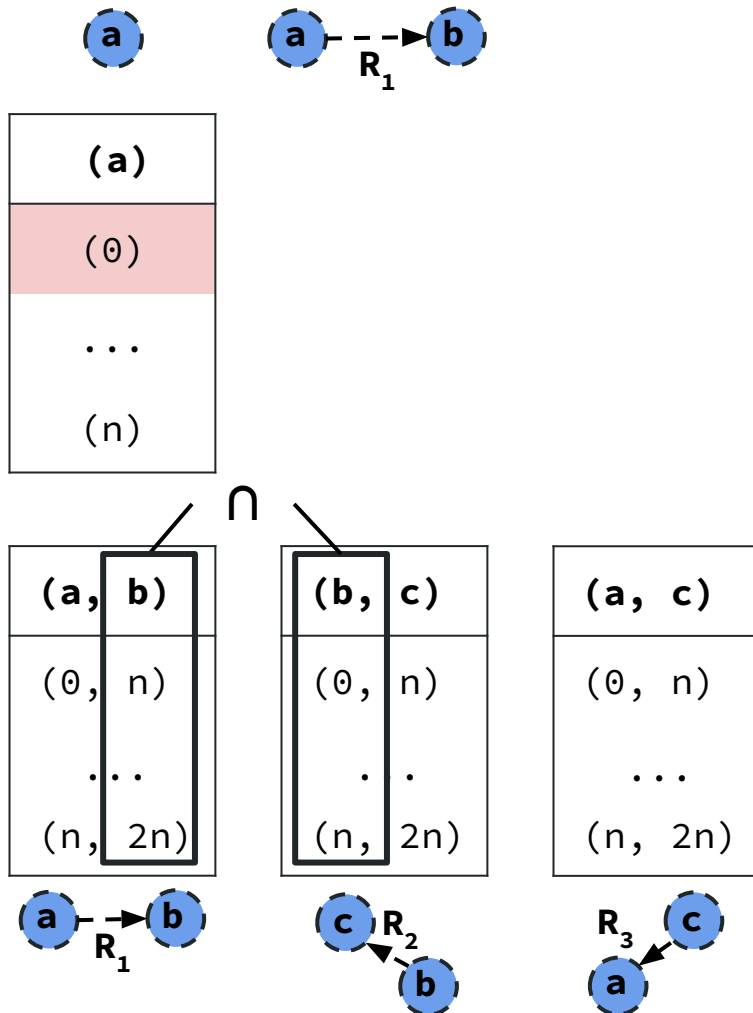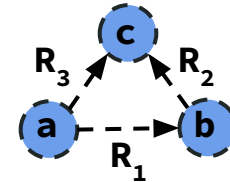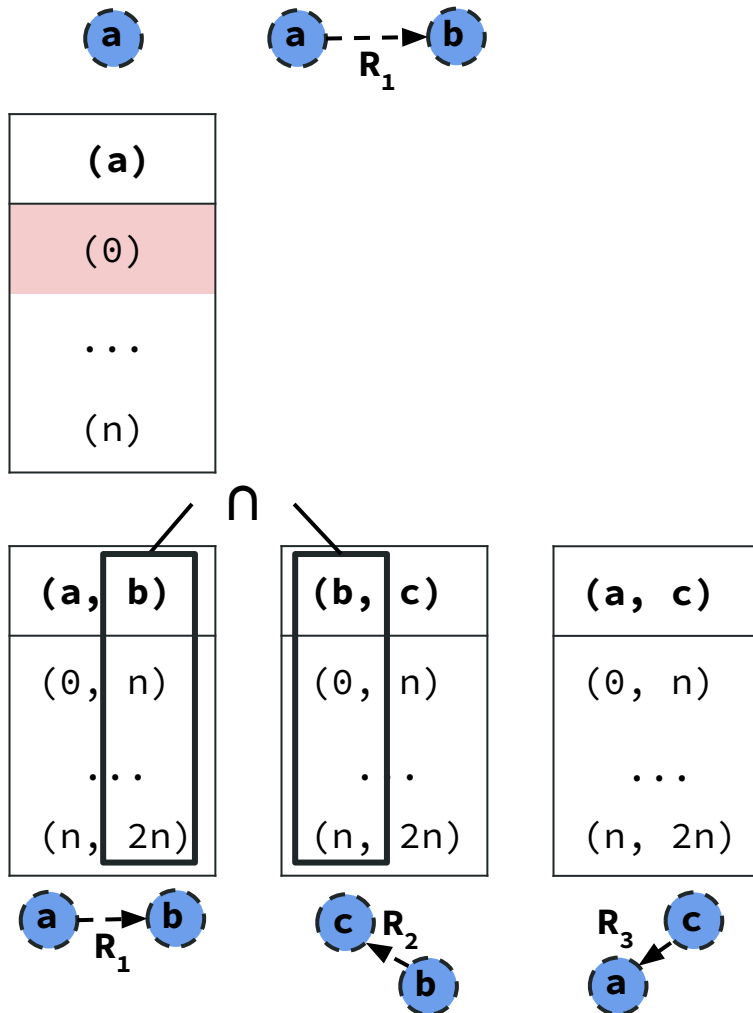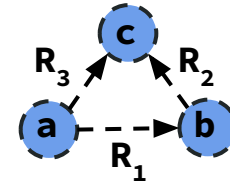
Query
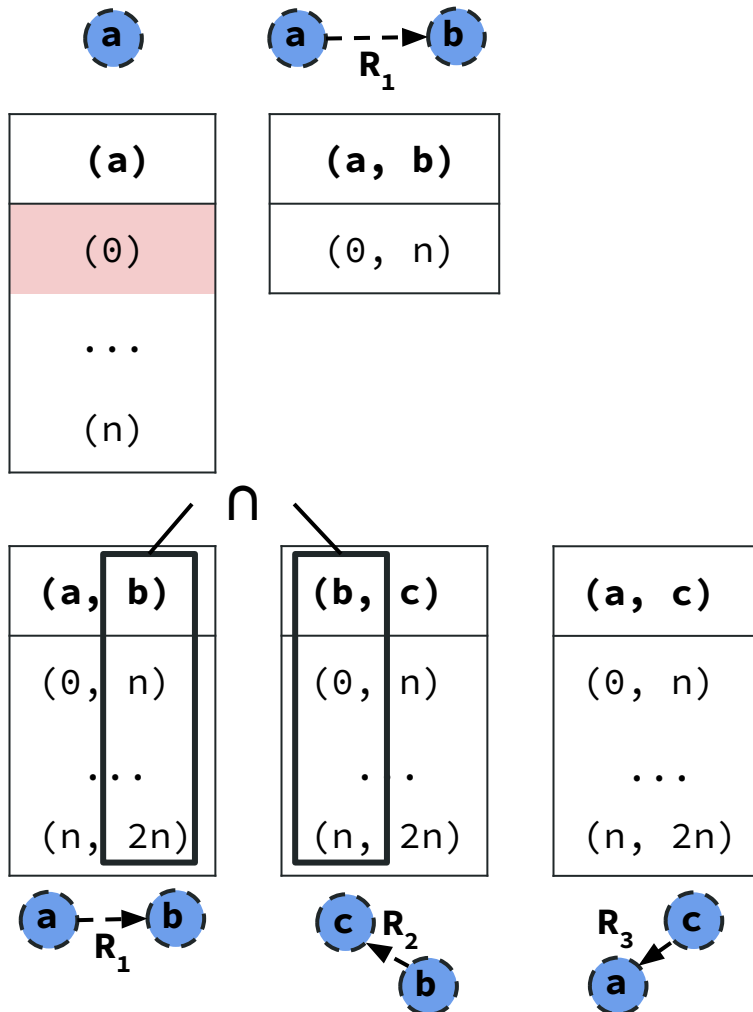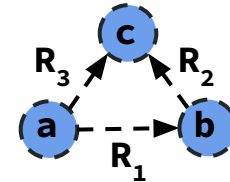

Dataset

(a, b)	⋈	(b, c)	⋈	(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Join Attribute Ordering (JAO)

[a,b,c]

Query

Dataset

$\cap$

(a, b)
(0, n)
...
(n, 2n)

a --> b
R_1

(b, c)
(0, n)
...
(n, 2n)

c R_2
b

(a, c)
(0, n)
...
(n, 2n)

R_3 c
a

92

Join Attribute Ordering (JAO)

[a,b,c]

Query

Dataset

$\cap$

(a, b)
(0, n)
..
(n, 2n)

a → b
R$_1$

(b, c)
(0, n)
...
(n, 2n)

c R$_2$
b

(a, c)
(0, n)
..
(n, 2n)

R$_3$ c
a

93

Join Attribute Ordering (JAO)

[a,b,c]

Query

Dataset

(a)
(0)
...
(n)

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

94

Join Attribute Ordering (JAO)

[a,b,c]

Query

(a)
(0)
...
(n)

Dataset

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

95

Join Attribute Ordering (JAO)

[a,b,c]

Query

(a)
(0)
...
(n)

Dataset

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Join Attribute Ordering (JAO)

[a,b,c]

Query

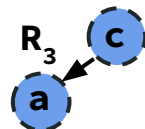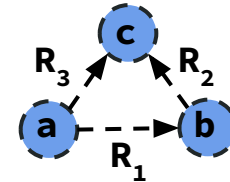| (a) |
|-----|
| (0) |
| ... |
| (n) |

Dataset

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

97

Join Attribute Ordering (JAO)

[a,b,c]

a

a $\dashrightarrow$ b
R_1

(a)
(0)
...
(n)

(a, b)
(0, n)
...
(n, 2n)

a $\rightarrow$ b
R_1

(b, c)
(0, n)
. .
(n, 2n)

c R_2
b

(a, c)
(0, n)
...
(n, 2n)

R_3 c
a

Query

c
R_3 R_2
a $\dashrightarrow$ b
R_1

Dataset

0 → n+1
1
n
n+2
...
n−1
2n

98

Join Attribute Ordering (JAO)

[a,b,c]

Query

(a)

(0)

...

(n)

∩

(a, b)		(b, c)		(a, c)
(0, n)		(0, n)		(0, n)
...		...		...
(n, 2n)		(n, 2n)		(n, 2n)

Dataset

Join Attribute Ordering (JAO)

[a,b,c]

Query

Dataset

∩

(a, b)
(0, n)
...
(n, 2n)

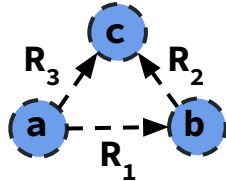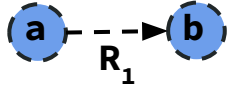| (b, c) |
| --- |
| (0, n) |
| .. |
| (n, 2n) |

(a, c)
(0, n)
...
(n, 2n)

(a)
(0)
...
(n)

R_1

R_1

R_2

R_3

R_3

R_2

R_1

Join Attribute Ordering (JAO)

[a,b,c]

Query

R_3 R_2

c

a $\xrightarrow{R_1}$ b

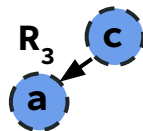| (a) |
|------|
| (0) |
| ... |
| (n) |

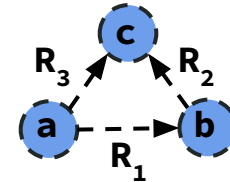| (a, b) |
|--------|
| (0, n) |

a $\xrightarrow{R_1}$ b

∩

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
..
(n, 2n)

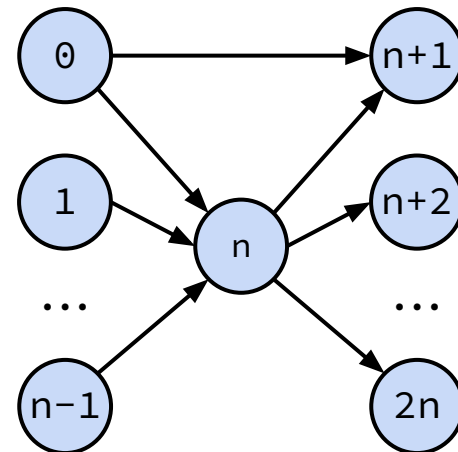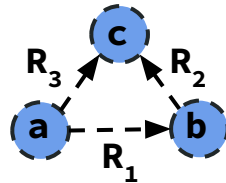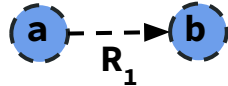| (a, c) |
|--------|
| (0, n) |
| ... |
| (n, 2n) |

a $\xrightarrow{R_1}$ b

c $\xrightarrow{R_2}$ b

R_3 c $\rightarrow$ a

Dataset

0 $\dashrightarrow$ n+1

1 $\rightarrow$ n $\rightarrow$ n+2

... n ...

n−1 $\rightarrow$ 2n

Join Attribute Ordering (JAO)

[a,b,c]

Query

(a)
(0)
...
(n)

(a, b)
(0, n)

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

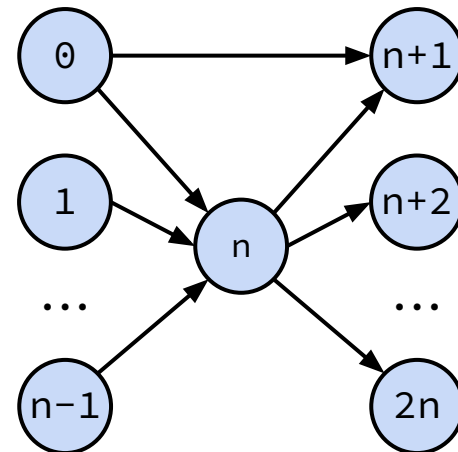| **(a, c)** |
|---|
| (0, n) |
| ... |
| (n, 2n) |

Dataset

102

Join Attribute Ordering (JAO)

[a,b,c]

(a)

(0)

...

(n)

(a, b)

(0, n)

(…, n)

(a, b)

(0, n)

...

(n, 2n)

(b, c)

(0, n)

...

(n, 2n)

(a, c)

(0, n)

...

(n, 2n)

Query

Dataset

Join Attribute Ordering (JAO)

[a,b,c]

a

a - - -> b
R_1

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

} N

(a, b)
(0, n)
...
(n, 2n)

a -> b
R_1

(b, c)
(0, n)
...
(n, 2n)

c R_2
b

(a, c)
(0, n)
...
(n, 2n)

R_3 c
a

Query

c

R_3 R_2

a - - -> b
R_1

Dataset

104

Join Attribute Ordering (JAO)

[a,b,c]

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

N

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Query

Dataset

105

Join Attribute Ordering (JAO)

[a,b,c]

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

$\}$ N

Query

$\cap$

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Dataset

106

Join Attribute Ordering (JAO)

[a,b,c]

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

(a, b, c)
(0, n, n+1)

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Query

Dataset

107

Join Attribute Ordering (JAO)

[a,b,c]

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

(a, b, c)
(0, n, n+1)

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Query

Dataset

Worst-case Optimal Joins (WCOJs) Example

Join Attribute Ordering (JAO)

[a,b,c]

WCOJs correct for sub-optimality
For all join attribute orderings.

→ *No advice on how to pick JAOs*

(a)
(0)
...
(n)

(a, b)
(0, n)
(…, n)

(a, b, c)
(0, n, n+1)

(a, b)
(0, n)
...
(n, 2n)

(b, c)
(0, n)
...
(n, 2n)

(a, c)
(0, n)
...
(n, 2n)

Dataset

Research Questions

Given an input Query,

1) How to pick good join attribute orderings?

Given an input Query,
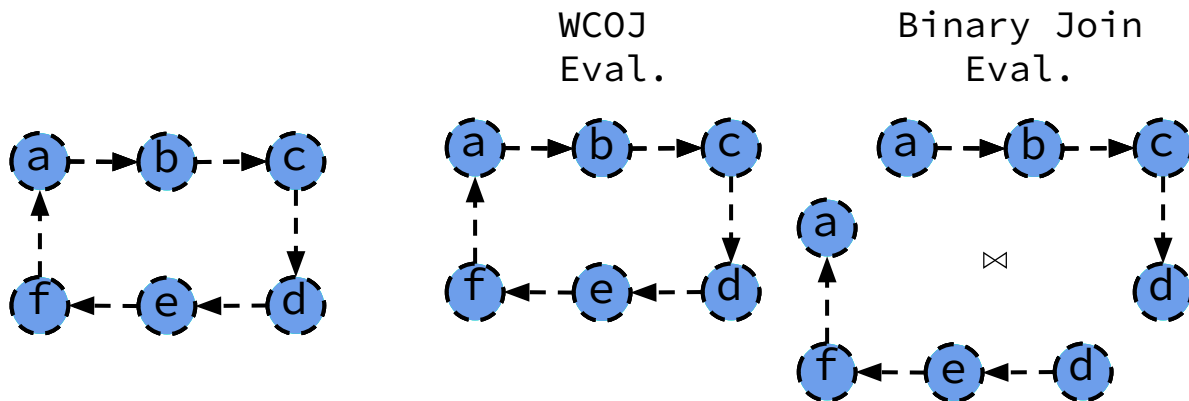
1) How to pick good join attribute orderings?

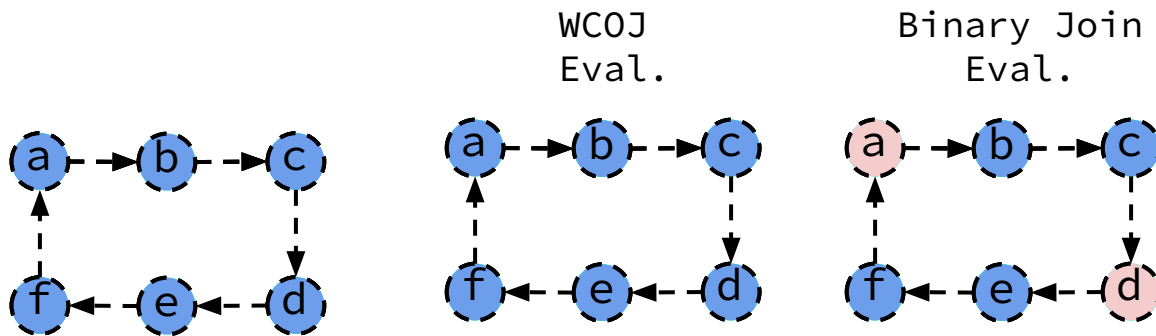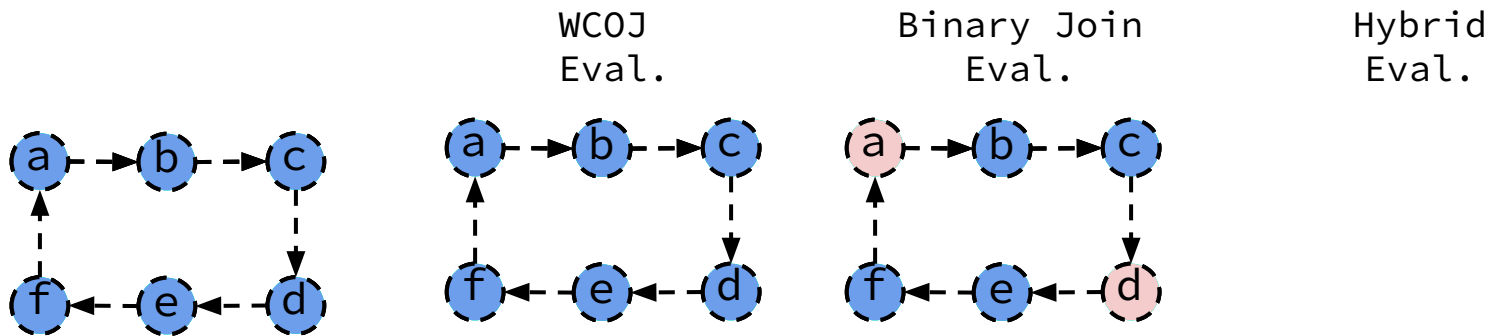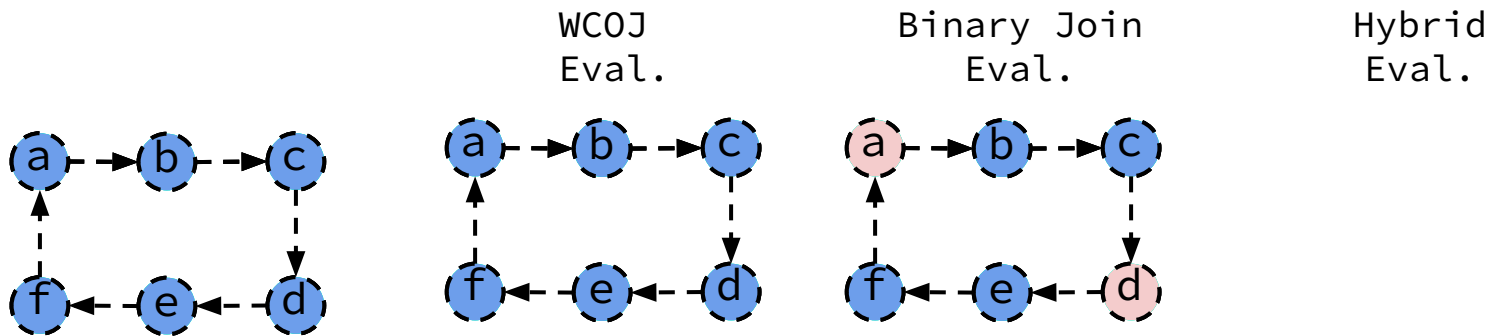2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,
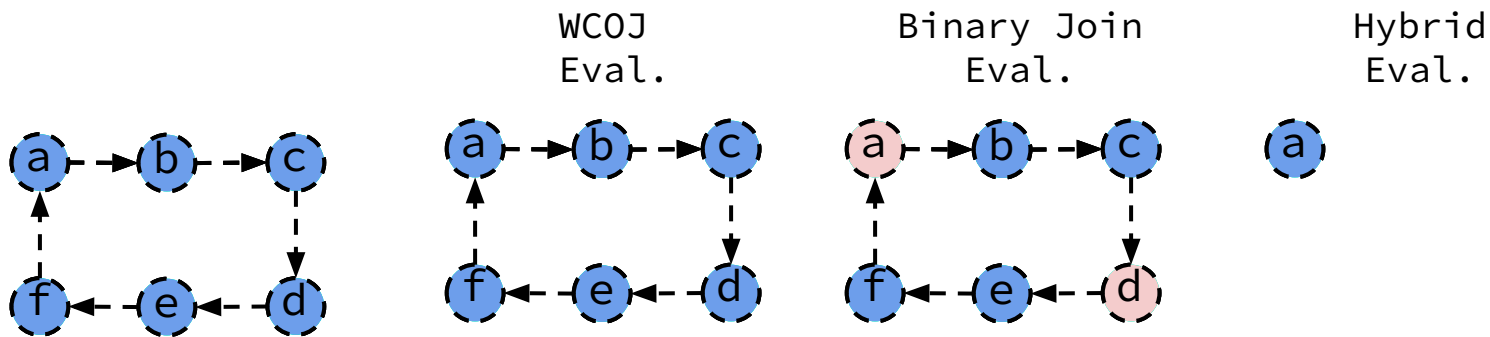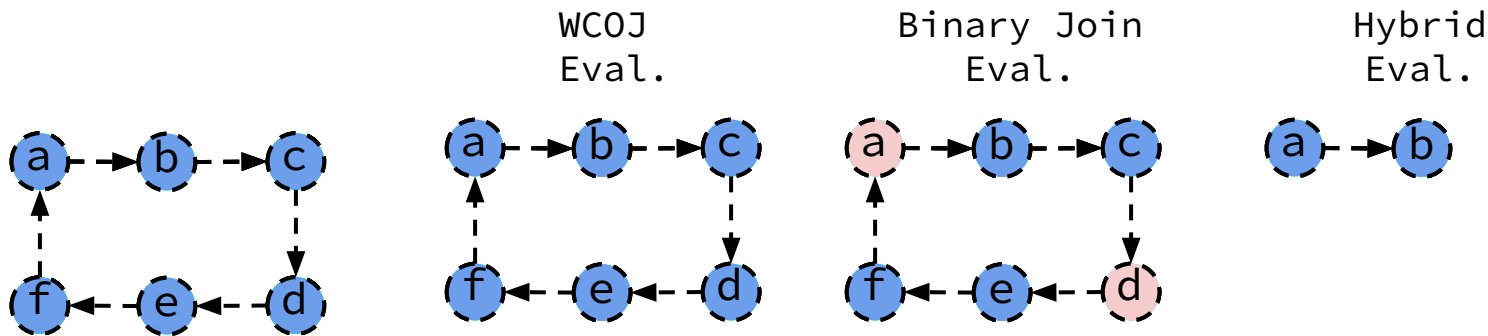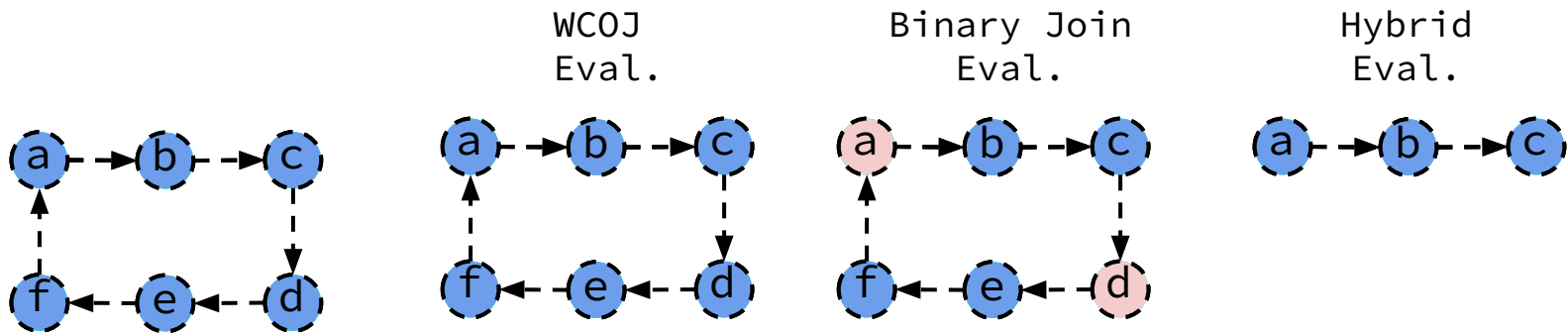
1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,
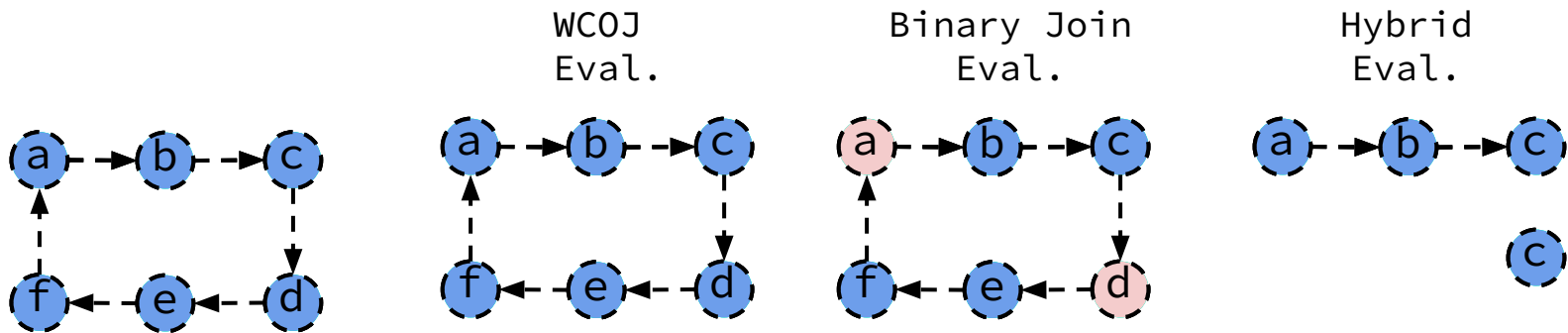
1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

```
WCOJ
Eval.
```

Given an input Query,
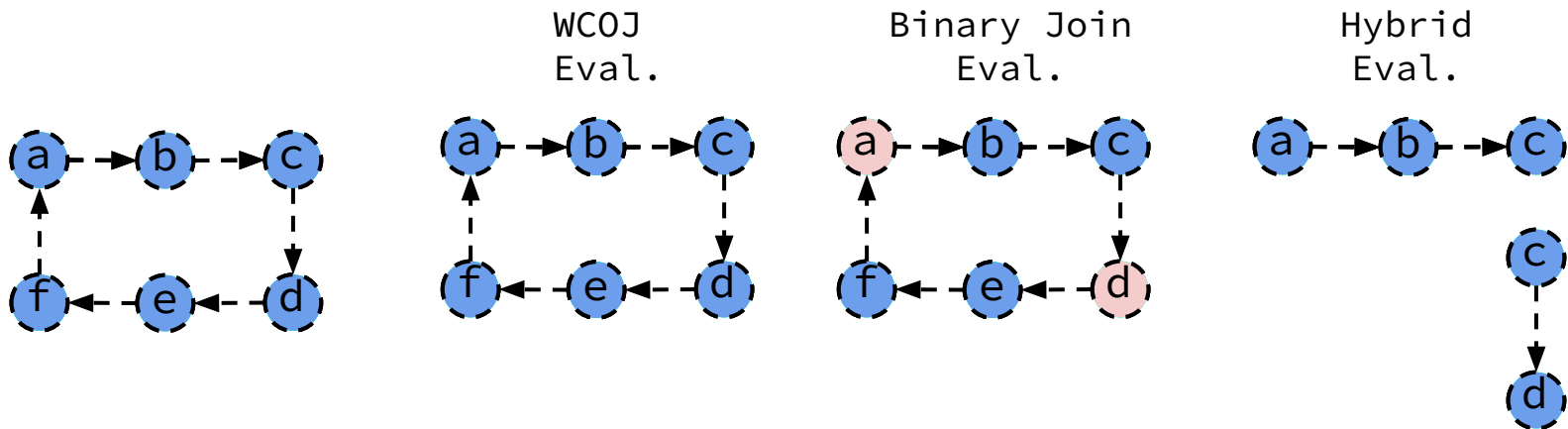
1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary
    joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary
   joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary
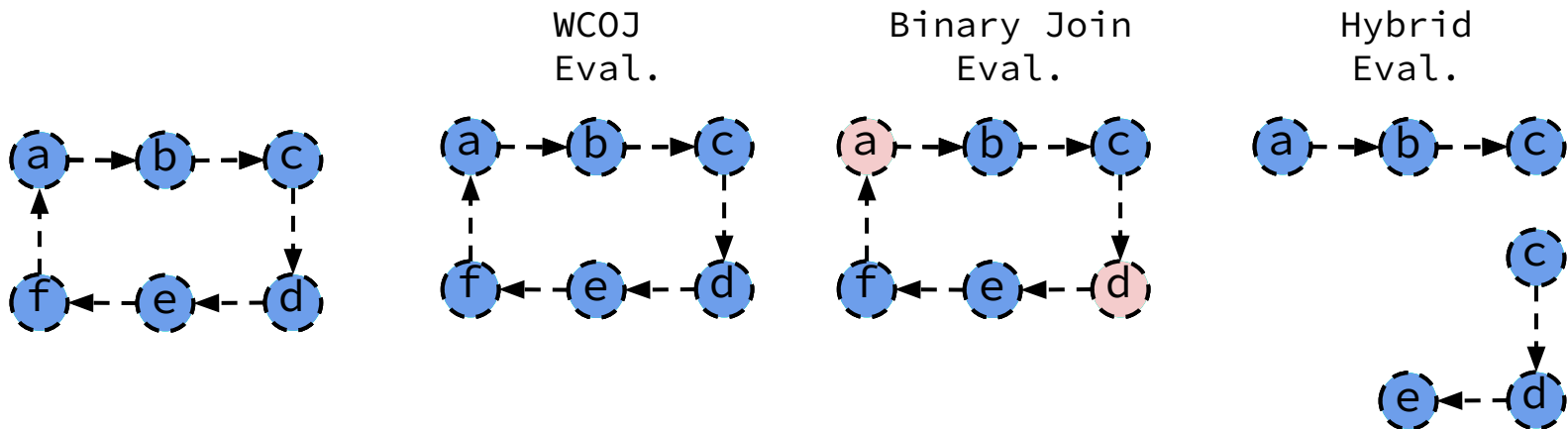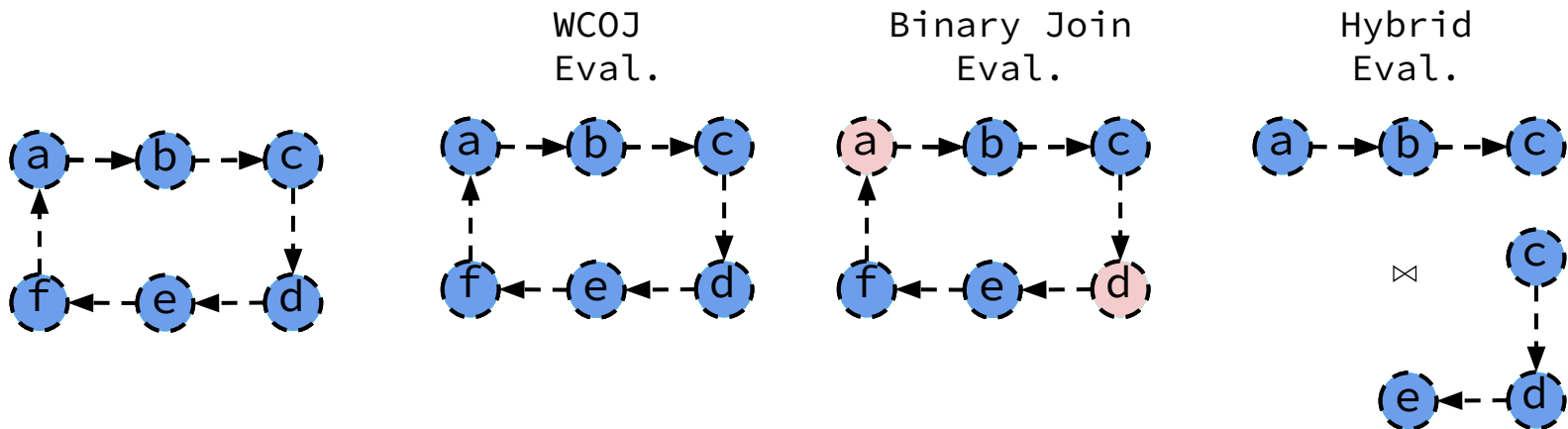   joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Given an input Query,
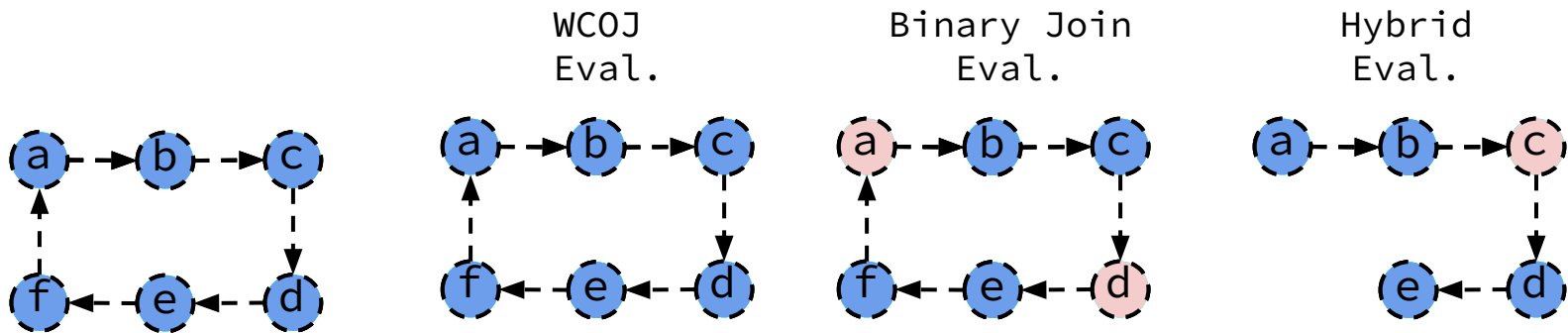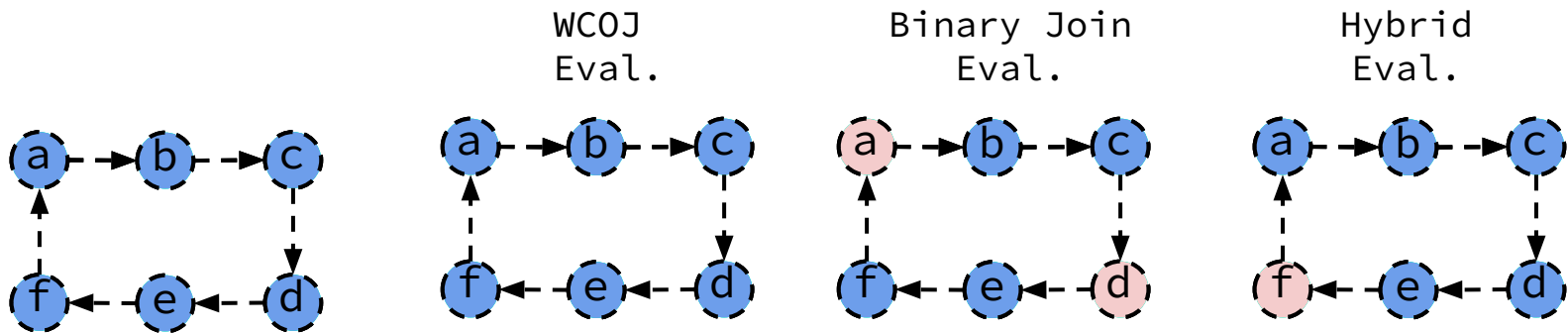
1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?
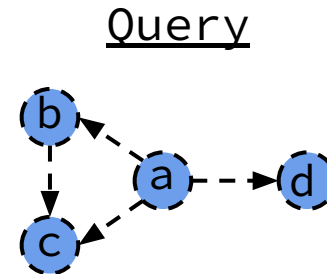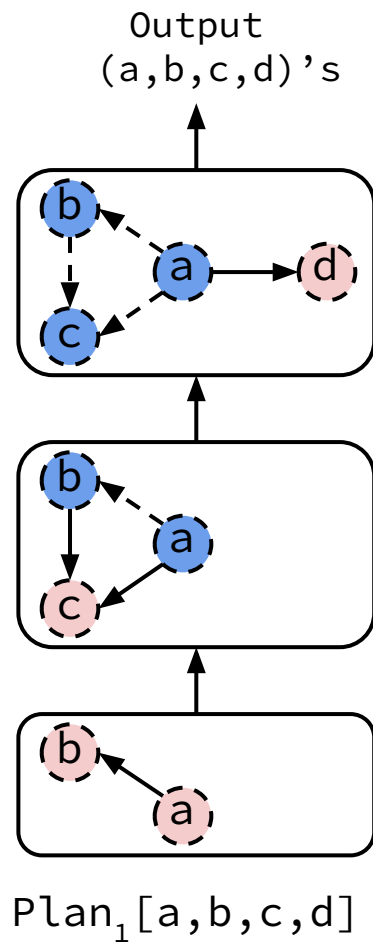

WCOJ
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

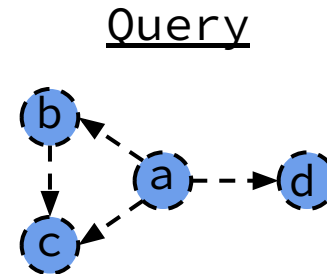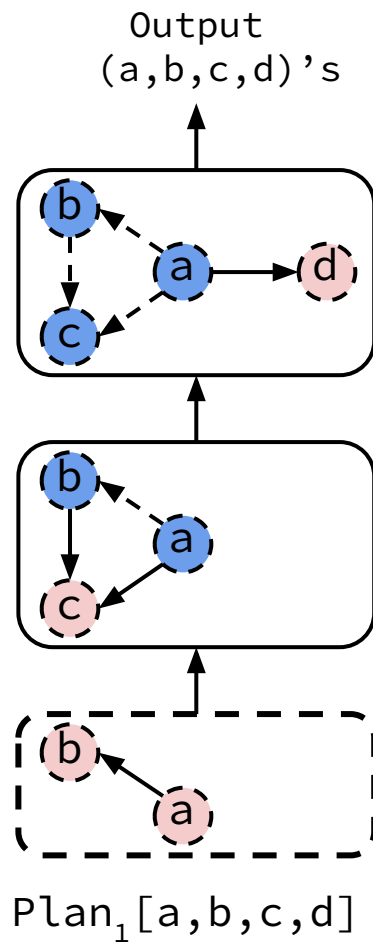Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

123

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

WCOJ
Eval.

Binary Join
Eval.

Hybrid
Eval.

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

Given an input Query,

1) How to pick good join attribute orderings?

2) How to generate efficient plans that mix binary joins and worst-case optimal joins (WCOJs)?

How to Pick Good Join Attribute Ordering (JAO)?

Query

Output
(a,b,c,d)'s

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan₁[a,b,c,d]

Query

Output
(a,b,c,d)'s

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

$(a,b) = (0,1)$

Scan (a,b)'s $= \{ (0,1), (0,2), (0,5), \ldots \}$

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

$(a,b)=(0,1)$
 c's: fwd$(a=0)$ ∩ fwd$(b=1)$
        $\{1,2,5,…\}$ ∩ $\{2,3,6,…\}$

Scan (a,b)'s = { (0,1),(0,2),(0,5),…}

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

$(a,b) = (0,1)$
 c's: fwd($a = 0$) ∩ fwd($b = 1$)
        {1, 2, 5, …} ∩ {2, 3, 6, …}

        { 2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
        {1, 2, 5, …} ∩ {2, 3, 6, …}

        { 2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
         {1, 2, 5, …} ∩ {2, 3, 6, …}
              ⬇
         { 2, 7, 9, …}
                    ⬇
Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)
d's: fwd(a = 0): {1, 2, 5, …}

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
        {1, 2, 5, …} ∩ {2, 3, 6, …}

        { 2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

(a,b,c,d)
(0, 1, 2, 1)

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)
d's: fwd(a = 0): {1, 2, 5, …}

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
        {1, 2, 5, …} ∩ {2, 3, 6, …}
        { 2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

WCOJ Evaluation Overview

(a,b,c,d)
(0, 1, 2, 1)
(0, 1, 2, 2)

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)
d's: fwd(a = 0): {1, 2, 5, …}

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
        {1, 2, 5, …} ∩ {2, 3, 6, …}

        {2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan₁[a,b,c,d]

(a,b,c,d)
(0, 1, 2, 1)
(0, 1, 2, 2)
(0, 1, 2, 5)

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)

d's: fwd(a = 0): {1, 2, 5, …}

(a,b) = (0,1)

c's: fwd(a = 0) ∩ fwd(b = 1)

{1, 2, 5, …} ∩ {2, 3, 6, …}

{2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

154

(a,b,c,d)
(0, 1, 2, 1)
(0, 1, 2, 2)
(0, 1, 2, 5)
...

Query

Output
(a,b,c,d)'s

(a,b,c) = (0,1,2)
d's: fwd(a=0): {1, 2, 5, …}

(a,b) = (0,1)
c's: fwd(a=0) ∩ fwd(b=1)
      {1, 2, 5, …} ∩ {2, 3, 6, …}

      {2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

$Plan_1[a,b,c,d]$

(a,b,c,d)

(a,b,c,d)
(0, 1, 2, 1)
(0, 1, 2, 2)
(0, 1, 2, 5)
...

Query

Output
(a,b,c,d)'s

(a,b) = (0,1)

c's: fwd(a = 0) ∩ fwd(b = 1)

{1, 2, 5, …} ∩ {2, 3, 6, …}

{2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

(a,b,c,d)
(0, 1, 2, 1)
(0, 1, 2, 2)
(0, 1, 2, 5)
...

Query

Output
(a,b,c,d)'s

(a,b) = (0,1)
 c's: fwd(a = 0) ∩ fwd(b = 1)
          {1, 2, 5, …} ∩ {2, 3, 6, …}

          { 2, 7, 9, …}

Scan (a,b)'s = { (0,1), (0,2), (0,5), …}

Plan$_1$[a,b,c,d]

```
Impact of JAO:
1) Number of intermediate results
2) Direction of Adj. Lists Intersected
```

How to Pick Good Join Attribute Ordering (JAO)?

Impact of JAO:

1) Number of intermediate results

2) Direction of Adj. Lists Intersected

Query

Output
(a,b,c,d)'s

Plan₁[a,b,c,d]

Query

(a,b,c,d)'s

(a,b,c,d)'s

Dataset	Plan	# Int. Results	Runtime (secs)
Epinions V=75K E=508K	Plan$_1$	4M	0.9
	Plan$_2$	55M **(13.8x)**	56.6 **(62.9x)**

Plan$_1$[a,b,c,d]

Plan$_2$[a,b,d,c]

How to Pick Good Join Attribute Ordering (JAO)?

Impact of JAO:

1) Number of intermediate results

2) Direction of Adj. Lists Intersected

Query

Query

Output
(a,b,c)'s

Plan$_1$
[a,b,c]

Query

(a,b,c)'s

(a,b,c)'s

Plan₁
[a,b,c]

Plan₂
[b,c,a]

Query

(a,b,c)'s

(a,b,c)'s

(a,b,c)'s

Plan$_1$
[a,b,c]

Plan$_2$
[b,c,a]

Plan$_3$
[a,c,b]

Query

(a,b,c)'s

(a,b,c)'s

(a,b,c)'s

Plan$_1$
[a,b,c]

Plan$_2$
[b,c,a]

Plan$_3$
[a,c,b]

Dataset	Plan	Σ Adj List sizes	Runtime (secs)
BerkStan V=685K E=7.6M	Plan$_1$	0.5B	2.6
	Plan$_2$	55B **(113.8x)**	15.2 **(5.8x)**
	Plan$_3$	55B **(114.0x)**	31.6 **(12.2x)**

How to Pick Good Join Attribute Ordering (JAO)?

Impact of JAO:

1) Number of intermediate results

2) Direction of Adj. Lists Intersected

→ Account for impact in **_cost model_**

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Output
(a,b,c)'s

c

a - - > b

a --> b

Plan$_1$
[a,b,c]

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider each tuple

Output
(a,b,c)'s

$$\sum_{t \in Q_{k-1}}$$

Plan$_1$
[a,b,c]

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider each tuple

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}}$$

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider each tuple

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}}$$

t = (a,b)
(0, 1)
(0, 2)
(1, 5)
...

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider each tuple's adj lists

Output
(a,b,c)'s

$$\sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

t = (a,b)
(0, 1)
(0, 2)
(1, 5)
...

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider the # input tuples

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

t = (a,b)
(0, 1)
(0, 2)
(1, 5)
...

|fwd(0)| + |fwd(1)|

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider the # input tuples

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

t = (a,b)	
(0, 1)	|fwd(0)| + |fwd(1)|
(0, 2)	|fwd(0)| + |fwd(2)|
(1, 5)	
...	

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider the # input tuples

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

t = (a,b)	
(0, 1)	|fwd(0)| + |fwd(1)|
(0, 2)	|fwd(0)| + |fwd(2)|
(1, 5)	|fwd(1)| + |fwd(5)|
...	

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Within an operator, consider the # input tuples

Output
(a,b,c)'s

Q_k

Q_{k-1}

Plan$_1$
[a,b,c]

$$\sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

t = (a,b)	
(0, 1)	|fwd(0)| + |fwd(1)|
	+
(0, 2)	|fwd(0)| + |fwd(2)|
	+
(1, 5)	|fwd(1)| + |fwd(5)|
	+
...	...

Operator I-cost

178

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Output
(a,b,c)'s

Plan$_1$
[a,b,c]

$$\sum_{Q_{k-1} \in Q_2 \ldots Q_{m-1}} \sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. } (i, dir) \text{ is accessed}}} |t[i].dir|$$

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Output
(a,b,c)'s

$$\sum_{Q_{k-1}\in Q_2...Q_{m-1}} \sum_{t\in Q_{k-1}} \sum_{\substack{(i,dir)\in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

1) number of
intermediate results

Plan$_1$
[a,b,c]

Cost Metric - Intersection Cost

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Output
(a,b,c)'s

$$\sum_{Q_{k-1}\in Q_2\ldots Q_{m-1}} \sum_{t\in Q_{k-1}} \sum_{\substack{(i,dir)\in A_{k-1} \\ \text{s.t. (i, dir) is accessed}}} |t[i].dir|$$

1) number of
intermediate results

2) size of adj. lists
dir. of adj. lists

Plan₁
[a,b,c]

Cost of a WCOJ plan is total **intersection-cost** of all operators.
I-cost: size of intersected adj lists throughout execution.

Output
(a,b,c)'s

Plan₁
[a,b,c]

$$\sum_{Q_{k-1} \in Q_2 \ldots Q_{m-1}} \sum_{t \in Q_{k-1}} \sum_{\substack{(i,dir) \in A_{k-1} \\ \text{s.t. } (i, dir) \text{ is accessed}}} |t[i].dir|$$

1) number of
intermediate results

2) size of adj. lists
dir. of adj. lists

I-Cost captures both effects!
Estimated using a sampling based approach.

182

Consider 2-join attributes

Consider 3-join attributes

Consider 4-join attributes

At each level:
- Consider the next Join Attribute
- Consider binary joins

Consider 4-join attributes

At each level:
- ● Consider the next Join Attribute
- ● Consider binary joins

Consider 4-join attributes

Mixing Binary and Worst-case Optimal Joins

At each level:
- Consider the next Join Attribute
- Consider binary joins

Consider 4-join attributes

At each level:
- Consider the next Join Attribute
- Consider binary joins

Consider 4-join attributes

Dynamic Programming Query & Dataset Cost-Based Optimizer

Graph_flow_

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

WCOJ
Subplan

Binary
Join

WCOJ
Subplan

Graph_flow_

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

EMPTYHEADED

**Generalized Hypertree Decomposition
Query Cost-Based Optimizer**

Dynamic Programming Query & Dataset Cost-Based Optimizer

Generalized Hypertree Decomposition Query Cost-Based Optimizer

This is leads to a specific join plan.

$$Q = Q_A \cup Q_B \cup Q_C \cup Q_D$$

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

**Generalized Hypertree Decomposition
Query Cost-Based Optimizer**

EMPTYHEADED

Generates only WCOJ subplans
followed by multiple binary joins

Example Graphflow Hybrid Plan

Dynamic Programming Query & Dataset Cost-Based Optimizer

Generalized Hypertree Decomposition Query Cost-Based Optimizer

Cost: maximum AGM bound of any of the leaves.
→ Minimize cost!

Example Graphflow Hybrid Plan

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

**Generalized Hypertree Decomposition
Query Cost-Based Optimizer**

Binary
Join

WCOJ
Subplan

Example Graphflow Hybrid Plan

Graphflow

**Dynamic Programming
Query & Dataset
Cost-Based Optimizer**

EMPTYHEADED

**Generalized Hypertree Decomposition
Query Cost-Based Optimizer**

Binary
Join

WCOJ
Subplan

Amazon V=403K E=3.4M	**Graphflow**	24.7 secs
	EmptyHeaded (EH)	> 30 mins
	EH in Graphflow	5.8 mins (**14x**)

Graphflow and EmptyHeaded Plan Spaces

Subgraph Queries: 14 queries

Dataset Domains: social networks, web, product co-purchasing

Differ in several structural properties:

(1) size
(2) how skewed their adjacency lists distribution is
(3) average clustering coefficients

Graphflow and EmptyHeaded Comparison

Subgraph Queries: 14 queries

Dataset Domains: social networks, web, product co-purchasing

# Joins	1-3	4-6	7+
+ Best JAO orderings			

Graphflow and EmptyHeaded Comparison

Subgraph Queries: 14 queries

Dataset Domains: social networks, web, product co-purchasing

# Joins	1–3	4–6	7+
Graphflow	–	1.8–3.2x Faster 10–25x in rare cases	Runs to completion 4-40 mins **72x speedup**
EMPTYHEADED + Best JAO orderings	1.5–2x faster	–	Timeout > 48 hrs

WCOJ Adoption

1. Plan Space

2. Cost Model (I-Cost)

3. Cardinality estimator

GraphScope

Alibaba

Outline

- **Novel Join Algorithms**
  → Worst-case Optimal Joins

- **Compressed Representations**
  → Factorized Representations

Flat Representation Example

Join Attribute Ordering (JAO): [b,a,c]

(b, a, c)

```
( b,   a,   c)
```

Flat Representation Example

```
( b,  a,  c)
```

(b, a, c)
(n, 0, n+1)
(n, 0, ...)
(n, 0, 2n)

n tuples

(b,   a,   c)
(n,   0, n+1)
(n,   0, ...)
(n,   0, 2n)
(n,   1, n+1)
(n,   1, ...)
(n,   1, 2n)

n tuples

n tuples

Flat Representation Example

(b,	a,	c)	
(n,	0,	n+1)	
(n,	0,	...)	n tuples
(n,	0,	2n)	
(n,	1,	n+1)	
(n,	1,	...)	n tuples
(n,	1,	2n)	
	...		
(n,	n−1,	n+1)	
(n,	n−1,	...)	n tuples
(n,	n−1,	2n)	

(b, a, c)		
(n, 0, n+1)		
(n, 0, ...)		
(n, 0, 2n)		
n, 1 n+1)		
n, 1 ...)		
n, 1 2n)		
...		
(n, n−1, n+1)		
(n, n−1, ...)		
(n, n−1, 2n)		

n tuples

n tuples

n tuples

Flat Representation Example

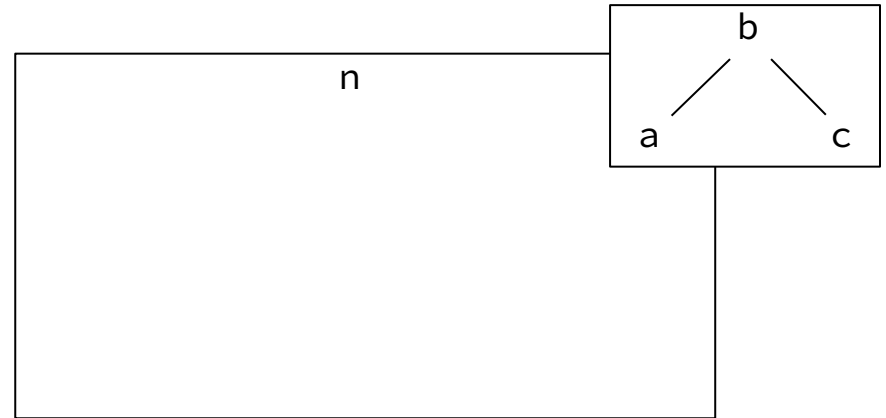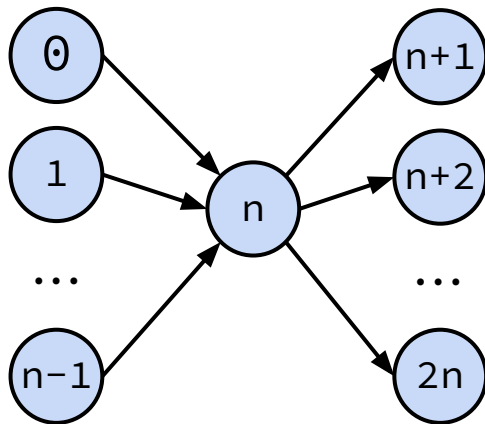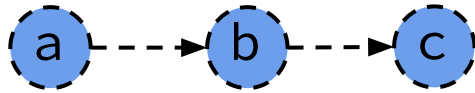| (b, a, c)		
(n, 0, n+1)		
(n, 0, ...)		n tuples
(n, 0, 2n)		
(n, 1, n+1)		
(n, 1, ...)		n tuples
(n, 1, 2n)		
...		
n, n−1 n+1)		
n, n−1 ...)		n tuples
n, n−1 2n)		

(b, a, c)
(n, 0, n+1)
(n, 0, ...)
(n, 0, 2n)
(n, 1, n+1)
(n, 1, ...)
(n, 1, 2n)
...
(n, n−1, n+1)
(n, n−1, ...)
(n, n−1, 2n)

n^2 tuples

a and c are
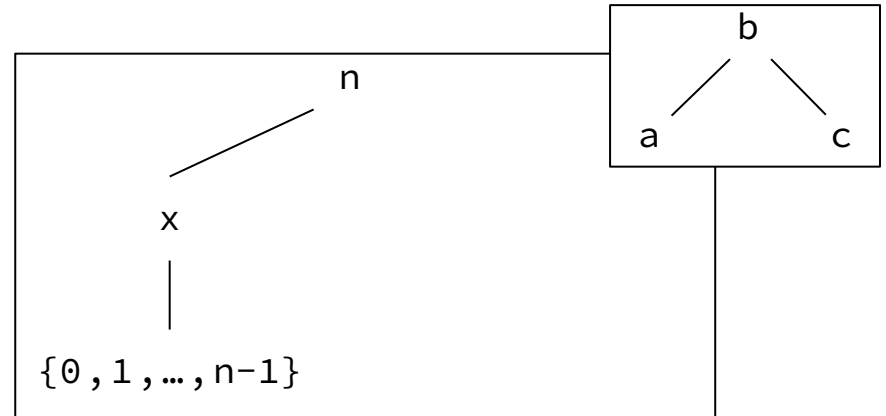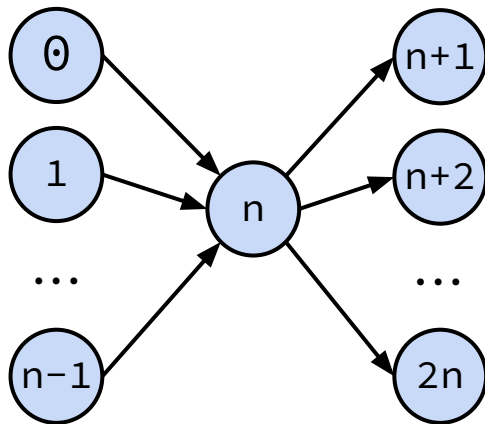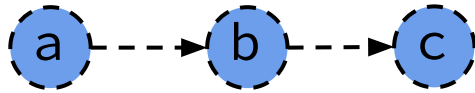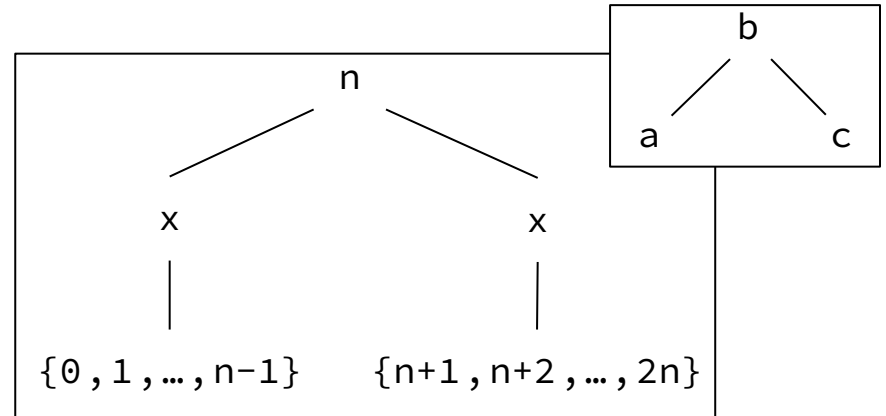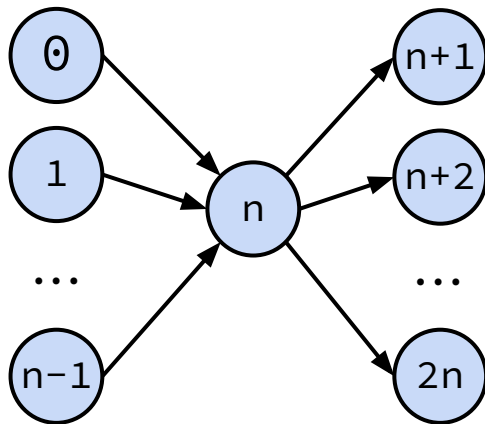*conditionally independent* on b!!

(b, a, c)
(n, 0, n+1)
(n, 0, ...)
(n, 0, 2n)
(n, 1, n+1)
(n, 1, ...)
(n, 1, 2n)
...
(n, n-1, n+1)
(n, n-1, ...)
(n, n-1, 2n)

n^2 **tuples**

2n+1 fields
vs.
3n² fields (flat)

Theory of factorization:

$$\sigma(Q) \leq AGM(Q)$$

where
$\sigma(Q)$: worst-case size bound
over f-representations.

In some cases,

$$\sigma(Q) < AGM(Q)$$

f-tree

f-representation

f-tree:
b
a   c

f-representation:
n
x   x
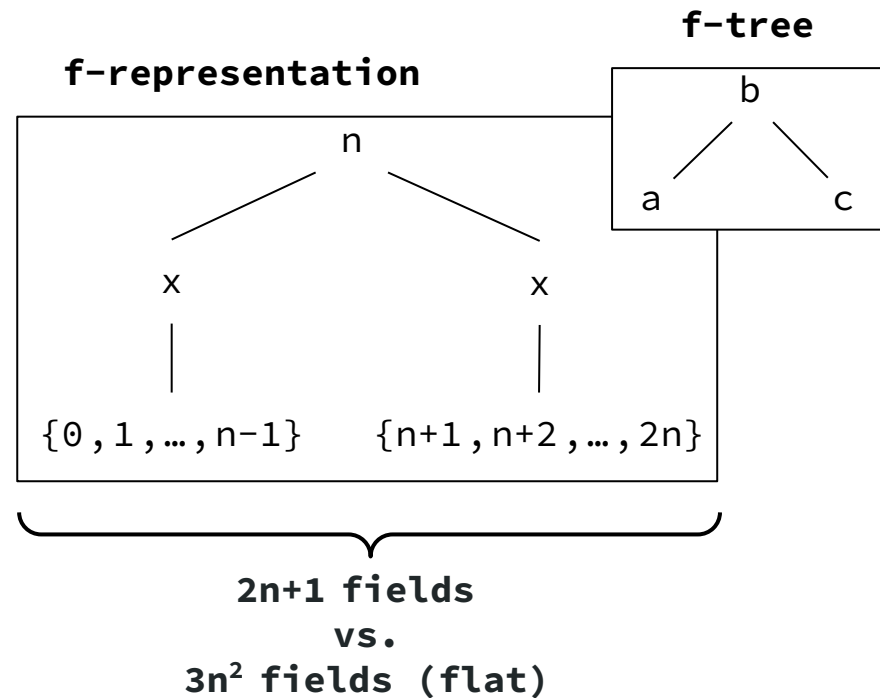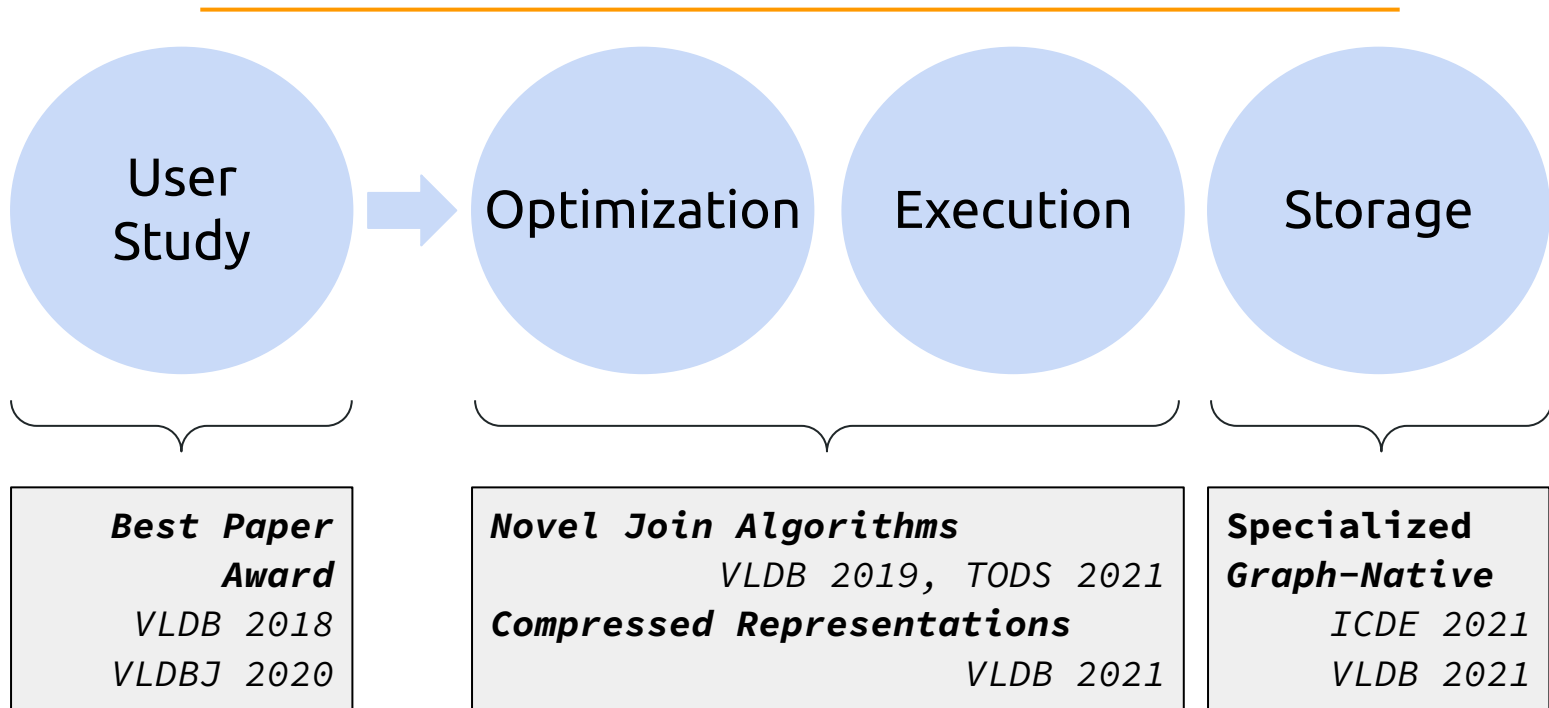{0,1,…,n-1}   {n+1,n+2,…,2n}

2n+1 fields
vs.
$3n^2$ fields (flat)

Fin.
Questions?

User
Study

Optimization

Execution

Storage

Best Paper Award
VLDB 2018
VLDBJ 2020

Novel Join Algorithms
VLDB 2019, TODS 2021
Compressed Representations
VLDB 2021

Specialized Graph-Native
ICDE 2021
VLDB 2021

Graph*flow*

SIGMOD 2017